

Special Topic: Cohesive Clustered Satellites System for 5GA and 6G Networks

# Game-theoretic clustering and scalable beamforming for multi-RIS-assisted cohesive satellite anti-jamming systems

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**Abstract** Cohesive clustered satellites have better scalability, spatial diversity, and coverage maintainability compared with conventional single satellite platforms. However, the high complexity of massive self-organizing satellites and malicious jamming of wireless channels pose severe threats to cohesive clustered satellite systems. By utilizing game-theoretic clustering to dynamically adjust the topology of cohesive clustered satellites, combined with the use of the reconfigurable intelligent surface (RIS) for reconfiguring the jamming environment, a multi-RIS-assisted cohesive clustered satellite anti-jamming system was developed for achieving reliable communication in the presence of malicious jamming. Specifically, game-theoretic clustering forms multiple coalitions for cohesive satellites, after which the satellite RIS can constructively enhance the desired signals while destructively weakening the jamming signals within one coalition. Building upon the above system, the aim of this study is to maximize the sum rate by jointly optimizing the satellite clustering, transceiver beamforming, and phase shifts of the satellite RIS when the jammer's channel state information (CSI) is imperfect. To solve the intractable problem, the satellite clustering was first optimized by formulating a coalition formation game, where a partial optimal coalition preference order was proposed for maximizing the sum rate. Subsequently, after converting the jammer's imperfect CSI into a robust CSI by adopting the discretization method, an alternative optimization method was developed for designing the transceiver beamforming and determining the phase shifts of the satellite RIS by leveraging the zero-forcing precoder, cyclic coordinate descent optimization framework, and linear minimum-mean-square-error decoder. The results prove that the proposed game-theoretic clustering method can converge to stable coalition formation with the exact potential game. The numerical simulation results demonstrate the superiority and validity of the proposed schemes compared to existing benchmark schemes.

**Keywords** satellite clustering, reconfigurable intelligent surface, coalition formation game, reliable communication, anti-jamming

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## 1 Introduction

Recently, cohesive clustered satellite communications have grown exponentially within the realms of space technology and information science [1–4]. Compared with conventional single satellites, cohesive clustered satellites have better scalability, spatial diversity, and coverage maintainability [5], thereby effectively improving the efficiency of space resource utilization. More specifically, through mutual cooperation, cohesive clustered satellites can achieve high-capacity communication with global coverage, thus enabling ubiquitous connectivity for remote areas and emergency situations.

However, the complexity of self-organizing networks between satellites increases with the number of satellites. To address this issue, satellite clustering has been adopted as a means of effectively improving the performance of cohesive satellite systems [6–12]. A satellite cluster framework that depends on the 3rd generation partnership project (3GPP) was proposed for improving the performance of cohesive satellite systems [13], where the efficiency and reliability of the entire system were enhanced by optimizing the communication protocols and standards among satellites. Moreover, a satellite routing algorithm with dynamic clustering, which effectively reduces the required storage space, was also proposed [14].

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Nevertheless, the above studies focused primarily on establishing clustering protocols or standards [15] that do not fully consider the topology of cohesive satellites.

Malicious jamming attacks also pose a severe threat to cohesive clustered satellites. To address this issue, various anti-jamming schemes have been widely adopted, such as the direct sequence spread spectrum and frequency-hopping [16–18]. However, both of these conventional schemes unavoidably consume additional bandwidth and energy. Although the multiple-input-multiple-output (MIMO) technique has been adopted for enhancing the desired signals and nullifying the jamming signals in low-earth-orbit (LEO) satellites [19], the implementation cost is prohibitive.

Fortunately, the reconfigurable intelligent surface (RIS), a newly emerging and innovative concept, offers an alternative scheme, in which the propagation environment can be reconfigured while maintaining low power consumption, thus boosting the desired signals and simultaneously weakening the jamming signals [20–26]. Specifically, RIS consists of massive low-cost passive reflective elements, which can dynamically change the phase shift of the desired signals, thereby improving communication performance and enhancing communication security [27]. An optimization framework was previously developed for the deployment of multi-RIS [28]. By employing graph-based network analysis, this study defines the RIS deployment problem and designs a sequential elimination method for cooperatively achieving the conflicting objectives of coverage maximization and cost minimization. Extensive simulations validated the superior cost-effectiveness of the graph-driven RIS deployment strategy compared to that of heuristic approaches. In a previous report [29], the performance of wireless energy transmission, especially the signal coverage and energy transmission efficiency in complex environments, was significantly improved by proposing and optimizing different beam-routing schemes. An innovative multi-RIS-aided downlink non-orthogonal multiple access (NOMA) framework was also designed for boosting the network throughput [30], where a sum-rate maximization problem was addressed through joint optimization of the reflection coefficients and power allocation. The simulation results demonstrated that the proposed RIS-NOMA architecture with optimized configurations attains substantial capacity improvements compared to conventional schemes.

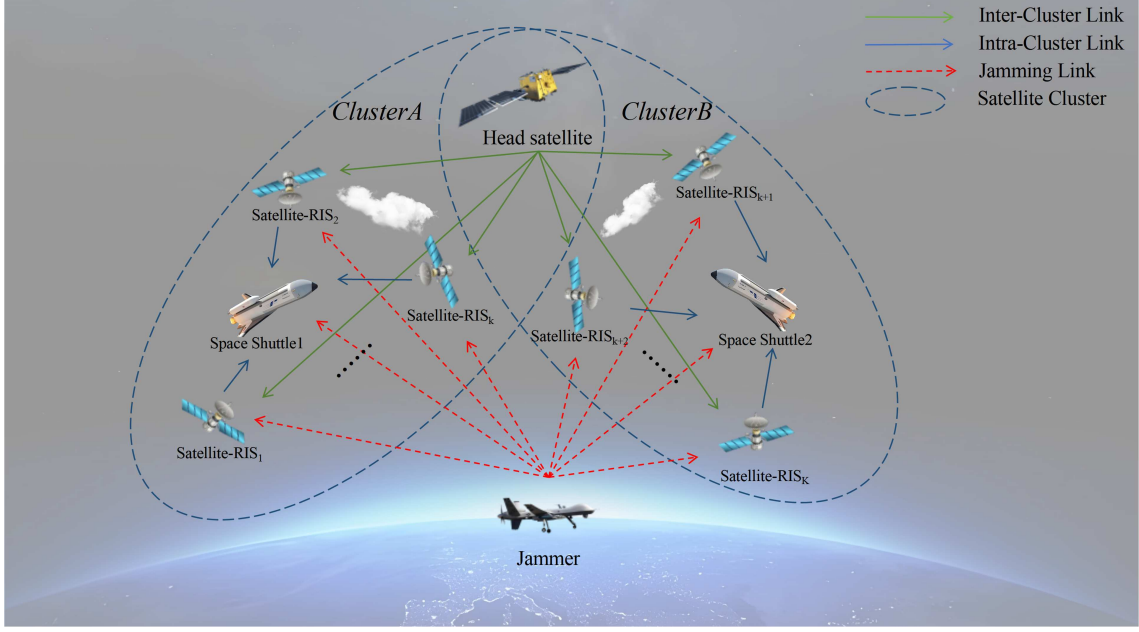
Thus, combined with the size, power, and weight constraints of satellites, RIS can be regarded as a promising solution for enabling cohesive clustered satellite systems [31–39]. An RIS-assisted satellite network architecture was previously proposed [31], in which RIS is deployed on the satellite to act as the reflective array for achieving significant transmission gain. In another study [32], an alternating optimization scheme was proposed for optimizing the satellite RIS reflection matrix, which significantly reduced the signal quality received by the eavesdropper. Moreover, a distortion-aware beamforming design for multi-beam satellite communications with nonlinear power amplifiers was also proposed [33], achieving good performance. However, the multi-RIS-assisted cohesive cluster satellite communication system has not been studied, especially for malicious jamming with imperfect channel state information (CSI).

Motivated by the above, this study presents a novel multi-RIS-assisted cohesive clustered satellite anti-jamming (CCSAJ) system. This kind of system can guarantee the reliability of communication when encountering malicious jamming. The main contributions of this study are as follows.

(1) To the best of our knowledge, this is the first report of a novel multi-RIS-assisted CCSAJ framework that defends against jamming attacks, where the game-theoretic satellite clustering can dynamically form multiple clustered satellite coalitions, and RIS can constructively enhance the desired signals, while destructively weakening the jamming signals within one coalition. In addition, in order to achieve the maximum sum rate, a worst-case optimization problem was formulated by jointly optimizing the satellite clustering, transceiver's beamforming and phase shifts of the satellite RIS under the condition of imperfect jamming CSI.

(2) To solve the intractable optimization problem, the problem is decoupled into two subproblems, i.e., optimization of satellite clustering and beamforming. For the design of satellite clustering, the progress of satellite clustering is first formulated as a coalition formation game (CFG), which can cope with the frequent and rapid changes in the coalition of the satellite. Thereafter, a partial coalition preference order is proposed for determining the optimal satellite clustering and maximizing the sum rate. The results demonstrate that the proposed CFG converges to a stable exact potential game (EPG) with at least one Nash equilibrium (NE).

(3) An alternative method is proposed for optimizing the transceiver's beamforming and phase shifts of the satellite RIS. To elaborate, a zero forcing (ZF) precoder is first adopted to optimize the transmit precoder, thereby eliminating inter-user interference. Subsequently, the cyclic coordinate descent (CCD) optimization algorithm is proposed to determine the closed-form solution of the phase shift for



**Figure 1** (Color online) Game-theoretic satellite clusters model.

multiple satellite RISs, leading to a significant decrease in computational complexity. Finally, the linear minimum-mean-square-error (MMSE) method is utilized to design the receive decoder. The simulation results clearly represent the effectiveness and superiority of the proposed schemes compared to existing approaches.

The layout of this paper is as follows. The system model and problem formulation are presented in Section 2. The proposed algorithm is presented in Section 3. Numerical results are provided in Section 4, and the conclusion is summarized in Section 5.

**Notation.**  $\mathbf{X}^H$ ,  $\mathbf{X}^T$ ,  $\mathbf{X}^*$ , and  $\|\mathbf{X}\|_F$  denote the conjugate transpose, transpose, conjugate, and Frobenius norm of the matrix  $\mathbf{X}$ . The notations  $\mathbb{E}\{\cdot\}$ ,  $\text{tr}\{\cdot\}$ ,  $\Re\{\cdot\}$ ,  $\Im\{\cdot\}$ , and  $\lambda\{\cdot\}$  denote the expectation, trace, real part, imaginary part, and eigenvalue of a complex number or matrix, respectively.  $\mathbb{C}^{m \times n}$  represents the complex space of  $m \times n$  dimensions. The symbol  $\mathbb{H}^{n \times n}$  is the Hermitian matrix of  $n \times n$  dimensions.  $[\cdot]_{n,n}$  represents the  $n$ -th diagonal element of the matrix. The distribution of a circularly symmetric complex Gaussian random vector with a mean vector  $x$  and covariance matrix  $\Sigma$  is denoted by  $\mathcal{CN}(x, \Sigma)$ .  $\text{diag}$  is a diagonal matrix.  $\mathcal{R}_+^K$  is the  $N$ -dimensional non-negative real domain.

## 2 System model and problem formulation

### 2.1 System model

As depicted in Figure 1, with the aid of  $n$  clustered satellite-RISs, a head satellite (HS, transmission satellite) with  $N_L$  antennas wishes to establish reliable links with  $k$  space shuttles. Meanwhile, a jammer utilizes a single omnidirectional antenna to impair the shuttles' signal reception from different angles. For simplicity of notion, the set of shuttles and satellite-RISs is denoted as  $\mathcal{K} = \{1, 2, \dots, K\}$  and  $\mathcal{N} = \{1, 2, \dots, N\}$ , respectively. We assume that each shuttle and satellite-RIS are equipped with  $N_U$  antennas and  $N_R$  RIS units, respectively. The channel coefficients of the HS-RIS <sub>$i$</sub>  link, the satellite-RIS <sub>$i$</sub> -shuttle <sub>$j$</sub>  link, the jammer-shuttle <sub>$j$</sub>  link and the jammer-satellite-RIS <sub>$i$</sub>  link are denoted by  $\mathbf{H}_{\text{HS}_i} \in \mathbb{C}^{N_R \times N_L}$ ,  $\mathbf{H}_{\text{RIS}_i \text{U}_j} \in \mathbb{C}^{N_U \times N_R}$ ,  $\mathbf{h}_{\text{JU}_j} \in \mathbb{C}^{N_U \times 1}$  and  $\mathbf{h}_{\text{JR}_i} \in \mathbb{C}^{N_R \times 1}$  respectively, which follow the channel model in [40]. Moreover, we assume that the reflection coefficients of the  $i$ -th satellite-RIS is  $\mathbf{p}_i = (p_1, \dots, p_{N_R})^T$ , where  $p_i = e^{j\theta_i}$  and  $\theta_i \in [0, 2\pi]$ . For ease of satellite, we denote the cascaded channels of HS-RIS <sub>$i$</sub> -shuttle <sub>$j$</sub>  link and jammer-RIS <sub>$i$</sub> -shuttle <sub>$j$</sub>  link as  $\mathbf{H}_{\text{HU}_{j,i}} = \mathbf{H}_{\text{RIS}_i \text{U}_j} \text{diag}(\mathbf{p}_i) \mathbf{H}_{\text{HS}_i}$  and  $\mathbf{h}_{\text{JU}_{j,i}} = \mathbf{H}_{\text{RIS}_i \text{U}_j} \text{diag}(\mathbf{p}_i) \mathbf{h}_{\text{JR}_i}$ , respectively.

Denote  $s_{T,j}$  as the symbol transmitted from HS to the  $j$ -th shuttle, where  $\mathbb{E}[|s_{T,j}|^2] = 1$ . Prior to transmission,  $s_{T,j}$  is weighted by the transmit beamforming vector  $\mathbf{w}_{T,j} \in \mathbb{C}^{N_L \times 1}$ . As such, the

signal transmitted to the  $j$ -th shuttle can be represented as  $\mathbf{x}_j = \mathbf{w}_{T,j} s_{T,j}$ . Considering the HS's power constraint, the transmit must satisfy  $\sum_{j=1}^K \|\mathbf{w}_{T,j}\|^2 \leq P_{\max}$ , where  $\|\cdot\|$  represents the Euclidean 2-norm, and  $P_{\max}$  is the maximum transmit power. Furthermore, the receive beamforming vector of the  $j$ -th shuttle is  $\mathbf{v}_j \in \mathbb{C}^{N_U \times 1}$ . Thus, the received signals at shuttle $_j$  is given by

$$y_j = \sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{T,j} s_{T,j} + \sum_{k \neq j} \sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{T,k} s_{T,k} + \sum_{i \in \mathcal{N}} \mathbf{v}_j^H (\mathbf{h}_{\text{JU}_{j,i}} + \mathbf{h}_{\text{JU}_j}) \mathbf{w}_{\text{JS}} s_{\text{J}} + \mathbf{v}_j^H \mathbf{n}_j, \quad \mathbf{n}_j \sim \mathcal{CN}(\mathbf{0}_{N_U}, \sigma_j^2 \mathbf{I}_{N_U}), \quad (1)$$

where  $\sigma_{U_j}^2$  is the white noise power. Accordingly, the shuttle $_j$ 's can be expressed as

$$\gamma_j = \frac{|\sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{T,j}|^2}{|\sum_{i \in \mathcal{N}} \mathbf{v}_j^H (\mathbf{h}_{\text{JU}_{j,i}} + \mathbf{h}_{\text{JU}_j}) \mathbf{w}_{\text{JS}}|^2 + |\sum_{j \neq k} \sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{T,k}|^2 + \sigma_j^2}. \quad (2)$$

In view of the non-cooperative relationship between the communication party and the jammer, it is a challenge to get the accurate CSI of the jammer. As a result, this paper models the jammer's CSI uncertainty as a given continuous angle range [40], which can be expressed as

$$\Delta = \{\mathbf{h}_{\text{JU}_j}, \mathbf{h}_{\text{JR}_i} | \theta \in [\theta_L, \theta_U], \varphi \in [\varphi_L, \varphi_U], \forall j, i\}, \quad (3)$$

where  $\theta, \varphi$  express the azimuth angle and elevation angle of the channels, respectively,  $(\theta_L, \theta_U)$  and  $(\varphi_L, \varphi_U)$  represent the bound of the corresponding angles.

It is worth noting that the transmission links in CCSAJ are modeled as quasi-static channels, ensuring that the temporal span of each transmission slot is confined within the channel coherence interval, thereby preserving time-selective channel stationarity. Consequently, the channel gain remains basically stationary in each transmission slot, but varies from different transmission slots [41].

## 2.2 Problem formulation

The objective of this paper is to improve transmission rate by jointly optimizing the satellite clustering, transceiver's beamforming, and satellite-RIS's phase shifts, under imperfect jammer's CSI. Accordingly, the problem can be modeled as

$$\begin{aligned} & \max_{\mathbf{w}_{T,j}, \mathbf{p}_i, \mathbf{v}_j, \text{Co}_j} \min_{\Delta} \sum_{j \in \mathcal{K}} R[\text{Co}_j], \\ & \text{s.t. C1: } \sum_{j=1}^K \|\mathbf{w}_{T,j}\|^2 \leq P_{\max}, \quad \text{C2: } |\mathbf{p}_i|_{n,n} = 1, \forall i, n, \\ & \text{C3: } \text{Co}_j \subseteq \mathcal{N}, \cup_{j \in \mathcal{K}} \text{Co}_j = \mathcal{N}, \text{Co}_i \cap \text{Co}_j = \emptyset, \forall i, j \in \mathcal{K}, i \neq j, \end{aligned} \quad (4)$$

where  $R[\text{Co}_j]$  is the sum rate of the cluster composed of the shuttle $_j$  and the serving satellite-RISs, i.e.,  $R[\text{Co}_j] = \log_2(1 + \gamma_j)$ , C1 is the transmitter power constraint,  $P_{\max}$  is the maximum transmitter power; C2 is the satellite-RIS's phase shift matrix amplitude constraint; C3 is the clustering constraint.

## 3 Game-theoretic satellite clustering and scalable beamforming

In this section, we first divide the problem into two subproblems: satellite clustering optimization, transceiver's beamforming design and satellite-RIS's phase shifts design. Then, two efficient algorithms for solving two subproblems are proposed, which are detailed in the following.

### 3.1 Game-theoretic clustering

In order to solve the optimization problem of satellite clustering, we use CFG for networking, mainly because CFG has significant advantages in autonomy, flexibility, incentive mechanism, resource optimization, scalability and robustness. It is suitable for a complex network environment, can effectively

improve network performance and stability, and is suitable for the scenarios considered in this article. In addition, compared with the traditional networking algorithm, CFG can realize the composition of an ad-hoc network according to the order, so as to maximize the utilization of resources.

**Definition 1** (CFG [42]). Each RIS can be arranged into a nonoverlapping coalitions  $\{Co_j\}$ ,  $\cup_{j \in \mathcal{K}} Co_j = \mathcal{N}$ , and each RIS has its own priority, which can be described as  $\mathcal{P} = (\succ_1, \succ_2, \dots, \succ_N)$ , then the strategy  $(\mathcal{N}, \mathcal{P})$  formulates CFG.

Each shuttle's utility function is defined as

$$u_i = R[Co_j \cup n_i] - R[Co_j], \forall n_i \notin Co_j. \quad (5)$$

Actually,  $u_i$  is the extra gain with the participation of  $n_i$ .

The coalition preference order is the guide of the coalition change principle, which influences the action of coalition member. The above two aspects determine the performance of CFG. Then, we provide the details in the following.

(1) Coalition preference order. In CFG, whether coalition members are willing to change the coalition is determined by the coalition preference order, which includes selfish order, Pareto order, and equilibrium order [43, 44].

**Definition 2** (Selfish order). If the coalition member  $n_i \in Co_{j'}$  and the coalition formations  $\{Co_j, Co_{j'}\}$ ,  $j, j' \in \mathcal{K}$  satisfy (6), the selfish order in CFG is formed, i.e.,

$$Co_j \succ_i Co_{j'} \Leftrightarrow R[Co_j \cup n_i] > R[Co_{j'} \setminus n_i]. \quad (6)$$

Following the selfish order, each coalition member prioritizes enhancing the sum rate of the newly formed coalition, disregarding the performance of the current coalition.

**Definition 3** (Pareto order). If the coalition member  $n_i \in Co_{j'}$  and the coalition formations  $\{Co_j, Co_{j'}\}$ ,  $j, j' \in \mathcal{K}$  satisfy (7), the Pareto order in CFG is formed, i.e.,

$$Co_j \succ_i Co_{j'} \Leftrightarrow R[Co_j \cup n_i] > R[Co_j] \& R[Co_{j'} \setminus n_i] > R[Co_{j'}]. \quad (7)$$

Different from the selfish order, the Pareto order has better stability due to the strong constraints, which make coalition members cannot leave the current cluster easily. It is shown that the selfish order focuses on the utility of individual coalition member, while the Pareto order pays more attention to the overall system. To balance the utility of individual and that of overall system, we propose a equilibrium order.

**Definition 4** (Equilibrium order). If the coalition member  $n_i \in Co_{j'}$  and the coalition formations  $\{Co_j, Co_{j'}\}$ ,  $j, j' \in \mathcal{K}$  satisfy (8), the equilibrium order in CFG is formed, i.e.,

$$\begin{aligned} Co_j \succ_i Co_{j'} &\Leftrightarrow u_i(Co_j \cup n_i) > u_i(Co_{j'}) \\ &\Leftrightarrow R[Co_j \cup n_i] - R[Co_j] > R[Co_{j'}] - R[Co_{j'} \setminus n_i] \\ &\Leftrightarrow R[Co_j \cup n_i] + R[Co_{j'} \setminus n_i] > R[Co_j] + R[Co_{j'}]. \end{aligned} \quad (8)$$

According to the equilibrium order, the coalition change of each coalition member will further improve the sum rate of both the target and current coalition, i.e.,  $Co_j$  and  $Co_{j'}$ , as well as the performance of the overall system.

(2) Coalition change principle. In this paper, we adopt the affiliation principle as the coalition change principle based on the proposed equilibrium order.

**Definition 5** (Affiliation principle). The coalition member  $n_i$  chooses to leave the current coalition  $Co_{j'}$  and enter the target coalition  $Co_j$ , if and only if the system get better sum rate according to the equilibrium order, namely,

$$n_i \rightarrow Co_j \Leftrightarrow u_i(Co_j \cup n_i) > u_i(Co_{j'}), \forall Co_j, Co_{j'} \quad \forall n_i \in Co_{j'}. \quad (9)$$

**Algorithm 1** Game-theoretic clustering algorithm.**input:** The set of satellite-RISs  $\mathcal{N}$  and shuttles  $\mathcal{K}$ .1: **Initialize:** Each satellite-RIS selects a shuttle to provide communication assistance randomly.2: **loop**3: Randomly select a satellite-RIS  $n_i$  and calculate its current coalition formation utility  $u_i(\text{Co}_{j'})$ ;4: The selected satellite-RIS  $n_i$  choose another shuttle  $j \in \mathcal{K}, j \neq j'$  for access and evaluate the explored utility  $u_i(\text{Co}_j)$ ;5: **if**  $u_i(\text{Co}_j) > u_i(\text{Co}_{j'})$  is satisfied **then**6: The satellite-RIS changes its current coalition and take part in the target one:  $\text{Co}_{j'} \rightarrow \text{Co}_{j'} \setminus n_i, \text{Co}_j \rightarrow \text{Co}_j \cup n_i$ ;7: **end if**8: **end loop:** All the satellite-RISs do not change their coalitions;**output:** Stable coalition formation  $\{\text{Co}_j\}^*$ .

**Definition 6** (Stable coalition formation). If any coalition member cannot improve the utility by changing its current coalition selection, the coalition formation is stable, which is given by

$$u_i(\text{Co}_j, \text{Co}_{-j}) \geq u_i(\text{Co}_{j'}, \text{Co}_{-j'}), \forall n_i \in \mathcal{N}, \text{Co}_j, \text{Co}_{j'} \in \{\cup_{j=1}^k \text{Co}_j\}. \quad (10)$$

**Theorem 1.** If equilibrium order is adopted as the coalition preference order, the proposed CFG can converge to a stable coalition formation.

*Proof.* EPG is employed to verify that a CFG with equilibrium order can reach a stable coalition formulation, which is defined as follows.

**Definition 7** (EPG [44]). If there is a potential function that satisfies (11), the game can be regarded as an EPG with at least one NE point:

$$\varphi(\text{Co}_j, \text{Co}_{-j}) - \varphi(\text{Co}_{j'}, \text{Co}_{-j'}) = u_i(\text{Co}_j) - u_i(\text{Co}_{j'}), \quad (11)$$

where  $\text{Co}_{-j}$  denotes the coalition without  $n_i$ .

We assumed the sum rate of the system as the potential function, i.e.,  $\varphi(\text{Co}_j, \text{Co}_{-j}) = R_{\text{sum}} = \sum_{j \in \mathcal{K}} R[\text{Co}_j]$ . When the coalition selection of satellite-RIS  $n_i$  changes from  $\text{Co}_{j'}$  to  $\text{Co}_j$ , the increase of potential function can be express as

$$\begin{aligned} \varphi(\text{Co}_j, \text{Co}_{-j}) - \varphi(\text{Co}_{j'}, \text{Co}_{-j'}) &= R[\text{Co}_j \cup n_i] + R[\text{Co}_{j'} \setminus n_i] \\ &\quad + R[\{\cup_{j \in \mathcal{K}} \text{Co}_j\} \setminus \{\text{Co}_j \cup n_i\} \setminus \{\text{Co}_{j'} \setminus n_i\}] \\ &\quad - R[\text{Co}_j] - R[\text{Co}_{j'}] - R[\{\cup_{j \in \mathcal{K}} \text{Co}_j\} \setminus \text{Co}_j \setminus \text{Co}_{j'}]. \end{aligned} \quad (12)$$

Since the change of the node's coalition strategy only impacts the achievable rate of the current coalition and the target coalition, which does not impact other coalitions, as such, we have

$$R[\{\cup_{j \in \mathcal{K}} \text{Co}_j\} \setminus \{\text{Co}_j \cup n_i\} \setminus \{\text{Co}_{j'} \setminus n_i\}] = R[\{\cup_{j \in \mathcal{K}} \text{Co}_j\} \setminus \text{Co}_j \setminus \text{Co}_{j'}]. \quad (13)$$

Combining (12) and (13), the increase of potential function can be further expressed as

$$\begin{aligned} \varphi(\text{Co}_j, \text{Co}_{-j}) - \varphi(\text{Co}_{j'}, \text{Co}_{-j'}) &= R[\text{Co}_j \cup n_i] + R[\text{Co}_{j'} \setminus n_i] - R[\text{Co}_j] - R[\text{Co}_{j'}] \\ &= R[\text{Co}_j \cup n_i] - R[\text{Co}_j] - (R[\text{Co}_{j'}] - R[\text{Co}_{j'} \setminus n_i]) \\ &= u_i(\text{Co}_j) - u_i(\text{Co}_{j'}). \end{aligned} \quad (14)$$

Observed from (11) and (14) can be found that the increase of the potential function and the coalition formation utility are equal. Thus, the CFG with equilibrium order is an EPG with at least one NE. Specifically, the potential function will increase when each satellite-RIS changes its coalition to improve the utility function. Moreover, due to the finite number of users and satellite-RISs, the possible results of coalition formulation are limited. After enough iterations, the coalition formulation with the highest potential function will be achieved. There must exist at least one coalition formulation to improve the utility if the final coalition formation  $\{\text{Co}_j\}^*$  is erratic. Finally, the stable coalition formulation is formed and the NE point is achieved. Theorem 1 is proved.

Finally, we provide the detailed game-theoretic clustering algorithm in Algorithm 1.

### 3.2 Scalable beamforming

After optimizing satellite clustering via CFG, we focus on designing the beamforming. Due to the imperfect jammer's CSI, we obtain robust CSI by discretizing the uncertainty region  $\Delta$ , and then applying the alternating optimization method to decouple and optimize multiple beamforming variables.



### 3.2.1 Discretization method for CSI uncertainty

To transform the intractable optimization problem in (4) with  $\Delta$  into a solvable one, we adopt the discretization method in [45] to address  $\Delta$ . In elaborate,  $\Delta$  is discretized into a set of  $S$  ( $S > 1$ ) samples, i.e.,

$$\tilde{\Delta} = \{\mathbf{h}_{\text{JU}_{j,1}}, \dots, \mathbf{h}_{\text{JU}_{j,S}}, \mathbf{h}_{\text{JR}_{i,1}}, \dots, \mathbf{h}_{\text{JR}_{i,S}}\}. \quad (15)$$

It is worth noting that as  $S \rightarrow \infty$ , discretizing  $\tilde{\Delta}$  approximates the continuity of  $\Delta$ . Furthermore, the azimuth angle and elevation angle inside  $\Delta$  are uniformly discretized, namely

$$\begin{aligned} \theta^{(p)} &= \theta_L + (k-1)\Delta\theta, k = 1, \dots, Q_1, \\ \varphi^{(q)} &= \varphi_L + (m-1)\Delta\varphi, m = 1, \dots, Q_2, \end{aligned} \quad (16)$$

where  $\theta^{(k)}$  and  $\varphi^{(m)}$  represent the azimuth angle and elevation angle of channel  $\mathbf{h}_{\text{JU}_j}^{(k,m)}$  and  $\mathbf{h}_{\text{JR}_i}^{(k,m)}$ , respectively,  $\Delta\theta = \frac{\theta_U - \theta_L}{Q_1 - 1}$ , and  $\Delta\varphi = \frac{\varphi_U - \varphi_L}{Q_2 - 1}$ . Thus, we can obtain the robust jammer's CSI  $\tilde{\mathbf{h}}_{\text{JU}_j} = \sum_{k=1}^{Q_1} \sum_{m=1}^{Q_2} (\frac{1}{Q_1 Q_2}) \mathbf{h}_{\text{JU}_j}^{(k,m)}$  and  $\tilde{\mathbf{h}}_{\text{JR}_i} = \sum_{k=1}^{Q_1} \sum_{m=1}^{Q_2} (\frac{1}{Q_1 Q_2}) \mathbf{h}_{\text{JR}_i}^{(k,m)}$ . By substituting the robust jammer's CSI into problem (4),  $\Delta$  can be removed. Therefore, the optimization problem can be reformulated as

$$\begin{aligned} \max_{\mathbf{w}_{\text{T},j}, \mathbf{P}_i, \mathbf{v}_j} \sum_{j \in \mathcal{K}} \log_2 & \left( 1 + \frac{|\sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{\text{T},j}|^2}{|\sum_{i \in \mathcal{N}} \mathbf{v}_j^H (\mathbf{h}_{\text{JU}_{j,i}} + \mathbf{h}_{\text{JU}_j}) \mathbf{w}_J|^2 + |\sum_{j \neq k} \sum_{i \in \mathcal{N}} \mathbf{v}_j^H \mathbf{H}_{\text{HU}_{j,i}} \mathbf{w}_{\text{T},k}|^2 + \sigma_j^2} \right), \\ \text{s.t. C1: } \sum_{j=1}^K \|\mathbf{w}_{\text{T},j}\|^2 & \leq P_{\max}, \quad \text{C2: } |[P_i]_{n,n}| = 1, \forall i, n, \end{aligned} \quad (17)$$

where C1 is the transmitter power constraint,  $P_{\max}$  is the maximum transmission power; C2 is the amplitude constraint of the RIS's phase shift matrix.

### 3.2.2 Beamforming design

First, we turn to optimizing the transmit beamforming  $\mathbf{w}_{\text{T},j}$ . Considering the system is a multi-user MIMO system, the ZF algorithm in [46] is the optimal solution for  $\mathbf{w}_{\text{T},j}$ . Specifically, the ZF algorithm transmits the signals towards the intended targets with nulls steered in the direction of other users, whose ZF matrix is

$$\mathbf{W}^\star = \beta_{\text{ZF}} \mathbf{G} (\mathbf{G}^H \mathbf{G})^{-1}, \quad (18)$$

where  $\beta_{\text{ZF}} = \sqrt{P_{\max} / \text{tr}((\mathbf{G}^H \mathbf{G})^{-1})}$ ,  $\mathbf{G}_j^H = \mathbf{v}_j^H (\sum_{i \in \text{Co}_j} \mathbf{H}_{\text{R}_i \text{U}_j} \text{diag}(\mathbf{p}_i) \mathbf{H}_{\text{LR}_i})$ , and  $\mathbf{G}_j^H$  is the  $k$ -th column of  $\mathbf{G}$ . Here,  $\mathbf{W}^\star = [\mathbf{w}_{\text{T},1}^\star, \dots, \mathbf{w}_{\text{T},K}^\star]$ , each column of which represents the transmit beamforming for the  $k$ -th shuttle.

After adopting the ZF algorithm to eliminate the inter-user interference, the beamforming subproblem can be reduced to a sum rate maximization problem within one cluster, i.e.,

$$\begin{aligned} \max_{\mathbf{p}_i, \mathbf{v}_j} \log_2 & \left( 1 + \frac{|\sum_{i \in \text{Co}_j} \mathbf{v}_j^H \mathbf{H}_{\text{LU}_{j,i}} \mathbf{w}_{\text{T},j}|^2}{|\sum_{i \in \text{Co}_j} \mathbf{v}_j^H (\mathbf{h}_{\text{JU}_{j,i}} + \mathbf{h}_{\text{JU}_j}) \mathbf{w}_J|^2 + \sigma^2} \right), \\ \text{s.t. C2: } |[p_i]_{n,n}| & = 1, \forall i, n. \end{aligned} \quad (19)$$

Due to the monotonicity of the logarithmic function, the objective function of (19) can also be reformulated as

$$\max_{\mathbf{p}_i, \mathbf{v}_j} \gamma = \frac{|\sum_{i \in \text{Co}_j} \mathbf{v}_j^H \mathbf{H}_{\text{LU}_{j,i}} \mathbf{w}_{\text{T},j}|^2}{|\sum_{i \in \text{Co}_j} \mathbf{v}_j^H (\mathbf{h}_{\text{JU}_{j,i}} + \mathbf{h}_{\text{JU}_j}) \mathbf{w}_J|^2 + \sigma^2},$$

$$\text{s.t. C2: } \left| [\mathbf{p}_i]_{n,n} \right| = 1, \forall i, n. \quad (20)$$

Define  $\mathbf{u}_{1,i}^H = \mathbf{v}_j^H \mathbf{H}_{R_i} \text{diag}(\mathbf{H}_{LR_i} \mathbf{w}_{T,j})$ ,  $\mathbf{u}_{2,i}^H = \mathbf{v}_j^H \mathbf{H}_{R_i} \text{diag}(\mathbf{h}_{JR_i} \mathbf{w}_J)$  and  $\hat{\mathbf{p}}_j = [\mathbf{p}_1, \dots, \mathbf{p}_i]$ ,  $i \in \text{Co}_j$  as intermediate variables. Removing the constant term in (20), the subproblem for optimizing the RIS's phase shifts is recast as

$$\begin{aligned} \max_{\hat{\mathbf{p}}_j} \quad & \gamma = \frac{|\hat{\mathbf{u}}_1^H \hat{\mathbf{p}}_j|^2}{|h_J + \hat{\mathbf{u}}_2^H \hat{\mathbf{p}}_j|^2 + \sigma^2}, \\ \text{s.t. C2: } \quad & \left| [\hat{\mathbf{p}}_j]_{n,n} \right| = 1, \forall n, \end{aligned} \quad (21)$$

where  $\hat{\mathbf{u}}_1^H \hat{\mathbf{p}}_j = \sum_{i \in \text{Co}_j} \mathbf{u}_{1,i}^H \mathbf{p}_i$ ,  $\hat{\mathbf{u}}_2^H \hat{\mathbf{p}}_j = \sum_{i \in \text{Co}_j} \mathbf{u}_{2,i}^H \mathbf{p}_i$ ,  $h_J = \sqrt{P_J} \mathbf{v}^H \mathbf{h}_{JU}$ . However, the objective function is a non-convex quadratic fractional function, which prevents us from solving it. To address this issue, the Dinkelbach method is adopted to convert the objective function of (21) into the following equivalent form, i.e.,

$$\hat{\mathbf{p}}_j^H \hat{\mathbf{U}}_1 \hat{\mathbf{p}}_j - \eta \left( |h_J|^2 + 2\Re(h_J \hat{\mathbf{u}}_2^H \hat{\mathbf{p}}_j) + \hat{\mathbf{p}}_j^H \hat{\mathbf{U}}_2 \hat{\mathbf{p}}_j + \sigma^2 \right), \quad (22)$$

where  $\eta$  is a non-negative Dinkelbach parameter,  $\hat{\mathbf{U}}_1 = \hat{\mathbf{u}}_1 \hat{\mathbf{u}}_1^H$ ,  $\hat{\mathbf{U}}_2 = \hat{\mathbf{u}}_2 \hat{\mathbf{u}}_2^H$ ,  $\mathbf{U} = \hat{\mathbf{U}}_1 - \eta \hat{\mathbf{U}}_2$ ,  $\hat{\mathbf{u}}_2^H = \eta h_J \hat{\mathbf{u}}_2^H$ . As such, we can obtain the optimization subproblem for  $\hat{\mathbf{p}}_j$ , namely

$$\begin{aligned} \max_{\hat{\mathbf{p}}_j} \quad & \hat{\mathbf{p}}_j^H \hat{\mathbf{U}} \hat{\mathbf{p}}_j - 2\Re(\tilde{\mathbf{u}}_2^H \hat{\mathbf{p}}_j), \\ \text{s.t. C2: } \quad & \left| [\hat{\mathbf{p}}_j]_{n,n} \right| = 1, \forall n. \end{aligned} \quad (23)$$

Then, we employ the CCD algorithm to solve it, which involves constructing a new vector  $[\eta; \hat{\mathbf{p}}_j]$  with  $\hat{N}_{R,j} + 1$  variables, where  $\hat{N}_{R,j} = \sum_{i \in \text{Co}_j} N_{R,i}$ . Specifically, we first expand the objective function of the equation, which is given by

$$\begin{aligned} \hat{\mathbf{p}}_j^H \hat{\mathbf{U}} \hat{\mathbf{p}}_j - 2\Re(\tilde{\mathbf{u}}_2^H \hat{\mathbf{p}}_j) &= \sum_{k=1}^{\hat{N}_{R,j}} \sum_{i=1}^{\hat{N}_{R,j}} \tilde{p}_{j,i}^* \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k} - 2\Re \left( \sum_{i=1}^{\hat{N}_{R,j}} \tilde{u}_{2,i}^* \tilde{p}_{j,k} \right) \\ &= \sum_{i=1}^{\hat{N}_{R,j}} \tilde{p}_{j,i}^* \hat{\mathbf{U}}_{(i,i)} \tilde{p}_{j,i} + \sum_{k \neq i}^{\hat{N}_{R,j}} \sum_{i=1}^{\hat{N}_{R,j}} \tilde{p}_{j,i}^* \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k} - 2\Re \left( \sum_{i=1}^{\hat{N}_{R,j}} \tilde{u}_{2,i}^* \tilde{p}_{j,k} \right) \\ &= \sum_{i=1}^{\hat{N}_{R,j}} \hat{\mathbf{U}}_{(i,i)} + \Re \left\{ \sum_{i=1}^{\hat{N}_{R,j}} \tilde{p}_{j,i}^* \left( \hat{\mathbf{U}}_{(i)} - 2\tilde{u}_{2,i} \right) \right\}, \end{aligned} \quad (24)$$

where  $\hat{\mathbf{U}}_{(i)} = \sum_{k=1}^{k < i} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k} + \sum_{k > i}^{\hat{N}_{R,j}} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k}$ . Armed with (24), a subproblem involving  $\hat{N}_{R,j} + 1$  variables can be obtained, and we update the above variables by using the CCD algorithm, which is expressed as

$$\left\{ \begin{aligned} \eta^{(i_d+1)} &= \frac{|\hat{c}_1^H \hat{\mathbf{p}}_j^{(i_d)}|^2}{|h_J + \hat{c}_2^H \hat{\mathbf{p}}_j^{(i_d)}|^2 + \sigma_U^2}, \\ \tilde{p}_{j,1}^{(i_d+1)} &= \arg \max_{\tilde{p}_{j,1} \in \mathcal{A}} e \left( \eta^{(i_d+1)}, \tilde{p}_{j,1}, \tilde{p}_{j,2}^{(i_d)}, \dots, \tilde{p}_{j,\hat{N}_{R,j}}^{(i_d)} \right), \\ &\vdots \\ \tilde{p}_{j,i}^{(i_d+1)} &= \arg \max_{\tilde{p}_{j,i} \in \mathcal{A}} e \left( \eta^{(i_d+1)}, \tilde{p}_{j,1}^{(i_d+1)}, \dots, \tilde{p}_{j,i-1}^{(i_d+1)}, \tilde{p}_{j,i}, \tilde{p}_{j,i+1}^{(i_d)}, \dots, \tilde{p}_{j,\hat{N}_{R,j}}^{(i_d)} \right), \\ &\vdots \\ \tilde{p}_{j,\hat{N}_{R,j}}^{(i_d+1)} &= \arg \max_{\tilde{p}_{j,\hat{N}_{R,j}} \in \mathcal{A}} e \left( \eta^{(i_d+1)}, \tilde{p}_{j,1}^{(i_d+1)}, \tilde{p}_{j,2}^{(i_d+1)}, \dots, \tilde{p}_{j,\hat{N}_{R,j}-1}^{(i_d+1)}, \tilde{p}_{j,\hat{N}_{R,j}} \right), \end{aligned} \right. \quad (25)$$



**Algorithm 2** Proposed CFG-BF optimization algorithm.**input:** The set of RISs  $\mathcal{N}$  and shuttles  $\mathcal{K}$  and the CSI of jamming channel.

1: **Initialize:** Set the iteration index  $t = 0$ . Generate the initial beams  $\{\{\mathbf{w}_{T,j}\}^{(0)}, \{\mathbf{p}_i\}^{(0)}, \{\mathbf{v}_j\}^{(0)}\}$  and the initial coalition  $\{\text{Co}_j\}^{(0)}$ ;  
2: **repeat**  
3:   For given  $\{\{\mathbf{w}_{T,j}\}^{(t)}, \{\mathbf{p}_i\}^{(t)}, \{\mathbf{v}_j\}^{(t)}\}$ , update  $\{\text{Co}_j\}^{(t)}$  to  $\{\text{Co}_j\}^{(t+1)}$  through Algorithm 1;  
4:   **if** coalition formation  $\{\text{Co}_j\}$  changed **then**  
5:     Optimize  $\{\mathbf{w}_{T,j}\}^{(t)}$  to  $\{\mathbf{w}_{T,j}\}^{(t+1)}$  with ZF algorithm based on the new coalition  $\{\text{Co}_j\}^{(t+1)}$  and  $\{\{\mathbf{p}_i\}^{(t)}, \{\mathbf{v}_j\}^{(t)}\}$ ;  
6:     Optimize  $\{\mathbf{p}_i\}^{(t)}$  to  $\{\mathbf{p}_i\}^{(t+1)}$  with CCD algorithm based on the new coalition  $\{\text{Co}_j\}^{(t+1)}$  and  $\{\{\mathbf{w}_{T,j}\}^{(t+1)}, \{\mathbf{v}_j\}^{(t)}\}$ ;  
7:     Optimize  $\{\mathbf{v}_j\}^{(t)}$  to  $\{\mathbf{v}_j\}^{(t+1)}$  with MMSE algorithm based on the new coalition  $\{\text{Co}_j\}^{(t+1)}$  and  $\{\{\mathbf{w}_{T,j}\}^{(t+1)}, \{\mathbf{p}_i\}^{(t+1)}\}$ ;  
8:   **end if**  
9:    $t = t + 1$ ;  
10: **until** The coalition formulation is stable or  $t = t_{\max}$ ;  
**output:** Optimal cooperative coalition formation  $\{\text{Co}_j\}^*$  and scalable beamforming vector  $\{\{\mathbf{w}_{T,j}\}^{(*)}, \{\mathbf{p}_i\}^{(*)}, \{\mathbf{v}_j\}^{(*)}\}$ .

where  $\mathcal{A} = \{z | |z| = 1, z \in \mathbb{C}\}$ ,  $e(\cdot) = \Re\{\sum_{i=1}^{\hat{N}_{R,j}} \tilde{p}_{j,i}^* (\hat{\mathbf{U}}_{(i)} - 2\tilde{u}_{2,i})\}$ . As for each optimization variable, the optimization subproblem can be formulated as

$$\begin{aligned} \max_{\tilde{p}_{j,i}} \Re \left\{ \tilde{p}_{j,i}^* \left( \sum_{k=1}^{k < i} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k}^{(i_d+1)} + \sum_{k > i}^{\hat{N}_{R,j}} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k}^{(i_d)} - 2\tilde{u}_{2,i} \right) \right\}, \\ \text{s.t. C2: } |\tilde{p}_{j,i}| = 1, \forall i. \end{aligned} \quad (26)$$

Obviously, the optimal closed-form solution for the expression can be obtained, i.e.,

$$\tilde{p}_{j,i}^\star = \exp \left\{ j \arg \left( \sum_{k=1}^{k < i} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k}^{(i_d+1)} + \sum_{k > i}^{\hat{N}_{R,j}} \hat{\mathbf{U}}_{(i,k)} \tilde{p}_{j,k}^{(i_d)} - 2\tilde{u}_{2,i} \right) \right\}. \quad (27)$$

Alternately updating  $\eta$  and  $\tilde{\mathbf{p}}_j$ , we can get the optimal closed-form solution for RIS's phase shifts.

Finally, we optimize the receive decoder  $\mathbf{v}_j$ . According to [47], the linear MMSE decoder can be adopted to solve the problem. Thus, we can obtain the optimal closed-form solution for  $\mathbf{v}$ , whose expression is as below

$$\mathbf{v}_j^\star = \frac{\left( \bar{\mathbf{w}}_{T,j} \bar{\mathbf{w}}_{T,j}^H + P_J \bar{\mathbf{w}}_J \bar{\mathbf{w}}_J^H + \sigma_U^2 \mathbf{I}_{\hat{N}_R} \right)^\dagger \bar{\mathbf{w}}_{T,j}}{\left\| \left( \bar{\mathbf{w}}_{T,j} \bar{\mathbf{w}}_{T,j}^H + P_J \bar{\mathbf{w}}_J \bar{\mathbf{w}}_J^H + \sigma_U^2 \mathbf{I}_{\hat{N}_R} \right)^\dagger \bar{\mathbf{w}}_{T,j} \right\|}, \quad (28)$$

where  $\bar{\mathbf{w}}_{T,j} = \sum_{i=1}^n \mathbf{H}_{R_i U_j} \text{diag}(\mathbf{p}_i) \mathbf{H}_{H R_i} \mathbf{w}_{T,j}$ ,  $\bar{\mathbf{w}}_J = (\mathbf{h}_J + \sum_{i=1}^n \mathbf{H}_{R_i U_j} \text{diag}(\mathbf{p}_i) \mathbf{h}_{J R_i})$ .

### 3.3 Convergence and complexity analysis

We summarize the proposed CFG-BF algorithm in Algorithm 2 for a better understanding. Then, we prove the convergence of the proposed algorithm, which includes two parts, i.e., the game-theoretic clustering and the beamforming design. As stated in Theorem 1, the game-theoretic clustering tends to a stable coalition formation. Thus, this subsection proves the convergence of the beamforming design, which mainly depends on the CCD. The ascent of (22) can be expressed as

$$\begin{aligned} e\left(\eta^{(id)}, \tilde{p}_{j,1}^{(id)}, \tilde{p}_{j,2}^{(id)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id)}\right) &\stackrel{a}{\leq} e\left(\eta^{(id)}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id)}\right) \\ &\stackrel{b}{\leq} e\left(\eta^{(id+1)}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id)}\right) \\ &\stackrel{a}{\leq} e\left(\eta^{(id+1)}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id+1)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id)}\right) \\ &\stackrel{b}{\leq} e\left(\eta^{(id+2)}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id+1)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id)}\right) \\ &\vdots \\ &\stackrel{a}{\leq} e\left(\eta^{(id+\hat{N}_R-1)}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id+1)}, \dots, \tilde{p}_{j,\hat{N}_R}^{(id+1)}\right) \end{aligned}$$

$$\leq e \left( \eta^{(id+\hat{N}_{R,j})}, \tilde{p}_{j,1}^{(id+1)}, \tilde{p}_{j,2}^{(id+1)}, \dots, \tilde{p}_{j,\hat{N}_{R,j}}^{(id+1)} \right). \quad (29)$$

The inequality  $a$  holds from the fact that  $\tilde{p}_{j,i}^{(id)}$  is a unique element-wise maximizer of  $e(\eta, \tilde{p}_j)$ , which will be given next. The inequality  $b$  holds based on [48].

To further prove the proposed algorithm's convergence, we discuss the uniqueness of the element-wise maximizer of  $\hat{p}_{j,i}$ . Recalling the solution (27) of problem (26), namely,  $\hat{p}_{j,i} = \exp\{j \arg(\tilde{U}_{(i)})\}$ , where  $\tilde{U}_{(i)} = \sum_{k=1}^{k < i} \hat{U}_{(i,k)} \tilde{p}_{j,k}^{(id+1)} + \sum_{k > i}^{\hat{N}_{R,j}} \hat{U}_{(i,k)} \tilde{p}_{j,k}^{(id)} - 2\tilde{c}_{2,i}$ , we can show that Eq. (27) is the unique element-wise maximizer by contradiction. Denoting  $\tilde{p}_{j,i} = \exp\{j \arg(\tilde{U}_{(i)} + \theta_i)\}$  as another element-wise maximizer, where  $\theta_i$  is not an integer multiple of  $2\pi$ , we have

$$\begin{aligned} e(\eta, \tilde{p}_j) &= \sum_{i=1}^{\hat{N}_{R,j}} \hat{U}_{(i,i)} + \sum_{i=1}^{\hat{N}_{R,j}} |\tilde{U}_{(i)}|, \\ e(\eta, \tilde{p}_j) &= \sum_{i=1}^{\hat{N}_{R,j}} \hat{U}_{(i,i)} + \sum_{i=1}^{\hat{N}_{R,j}} |\tilde{U}_{(i)}| \cos(\theta_i). \end{aligned} \quad (30)$$

Since both  $\tilde{p}_j$  and  $\bar{p}_j$  maximize  $e(\eta, \tilde{p}_j)$ , we obtain  $e(\eta, \tilde{p}_j) = e(\eta, \bar{p}_j)$  so that  $\theta_i$  must be an integer multiple of  $2\pi$ , which contradicts our assumption. Thus, Eq. (27) is the unique element-wise maximizer of  $e(\eta, \tilde{p}_j)$ , such that inequality  $b$  holds, which implies that the sequence  $\{\eta^{(id)}, \tilde{p}_j^{(id)}\}$  is monotonically increasing. Furthermore, the objective function  $\{\eta^{(id)}, \tilde{p}_j^{(id)}\}$  is upper-bounded by the constraint C2. Hence,  $\{\eta^{(id)}, \tilde{p}_j^{(id)}\}$  increases to the limited point  $\{\eta^\infty, \tilde{p}_j^\infty\}$ . Thus, the proposed beamforming algorithm converges to the optimal solution. Combining the convergence of game-theoretic clustering and the beamforming design, the overall algorithm can converge to the fixed solution.

Then, we will prove the scalability of the proposed algorithm compared to existing algorithms by calculating its complexity. Specifically, the complexity of the proposed algorithm mainly includes three parts: satellite clustering based on CFG, CCD algorithm for (23) and optimization of  $\mathbf{w}_{T,j}$  and  $\mathbf{v}_j$ . We denote  $I_1$  and  $\varepsilon$  as the total number of alternating optimization algorithm's number and the fault tolerance of the CCD algorithm, respectively. According to [49], the complexity of updating a single  $\tilde{p}_i$  in (25) is  $\mathcal{O}(\hat{N}_{R,j})$ , and the complexity of CCD algorithm in (23) could be described as  $\mathcal{O}[(\frac{1}{\varepsilon^2})(\hat{N}_{R,j}^2 + 1)]$ . Next, we can get the complexity of optimizing  $\mathbf{w}_{T,j}$  and  $\mathbf{v}_j$  is  $\mathcal{O}(N_U \hat{N}_{R,j})$  and  $\mathcal{O}(N_T \hat{N}_{R,j})$  based on [7]. As such, the complexity of the alternating optimization algorithm can be expressed as

$$\mathcal{O}_{AO} = I_1 \left[ \mathcal{O} \left[ \left( \frac{1}{\varepsilon^2} \right) (\hat{N}_{R,j}^2 + 1) \right] + \mathcal{O}(N_U \hat{N}_{R,j}) + \mathcal{O}(N_T \hat{N}_{R,j}) \right]. \quad (31)$$

The proposed game-theoretic clustering algorithm involves the comparison between the current coalition and new coalition and the update for possible coalition switch where complexity can be represented as  $\mathcal{O}(\mathcal{U}_1)$  and  $\mathcal{O}(\mathcal{U}_2)$ , respectively. We denote  $I_2$  as the number of convergence iteration, and the complexity of CFG-based cohesive satellite clustering can be described as  $I_2 [\mathcal{O}(\mathcal{U}_1) + \mathcal{O}(\mathcal{U}_2)]$ . Thus, the total complexity could be expressed in the following form:

$$\mathcal{O}_{CCSAJ} = I_1 \left[ \mathcal{O} \left[ \left( \frac{1}{\varepsilon^2} \right) (\hat{N}_{R,j}^2 + 1) \right] + \mathcal{O}(N_U \hat{N}_{R,j}) + \mathcal{O}(N_T \hat{N}_{R,j}) \right] + I_2 [\mathcal{O}(\mathcal{U}_1) + \mathcal{O}(\mathcal{U}_2)]. \quad (32)$$

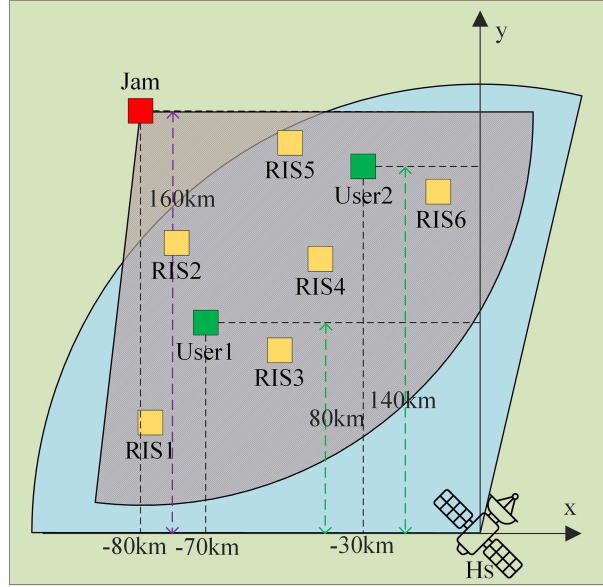
For the SCA algorithm, due to the optimization problem with  $\hat{N}_R + 3$  optimization variables and  $\hat{N}_R$  constraints in cvx. The complexity of SCA algorithm is  $\mathcal{O}[(\frac{1}{\varepsilon^2})\hat{N}_R(\hat{N}_{R,j} + 2)^2]$ . Compared with the complexity of the CCD algorithm, the proposed algorithm has better scalability.

## 4 Simulation and results

In this section, the performance of the proposed algorithm is demonstrated and verified through numerical simulation. The simulation parameters are summarized in Table 1. In this study, the jamming pattern of the hostile jammer is set to fixed power and frequency jamming, and the jamming power is set to

**Table 1** Simulation parameters.

Parameter	Symbol	Value
RIS number	$n$	6
Shuttle number	$k$	2
Transmission power	$P_{\max}$	10 W
Bandwidth	$B$	10 MHz
Noise power	$\sigma^2$	$10^{-11}$
Multi-path number	$N_m$	5
Jamming channel uncertainty	$\Delta$	$6^\circ$

**Figure 2** (Color online) Simulation scenario.

10 W, which is simplified. The RISs and shuttles are randomly deployed in the task area, as shown in Figure 2. Additionally, the path-loss factor  $g_d$  follows a complex Gaussian distribution with a mean of 0 and a variance of  $10^{\frac{PL}{10}}$ , where  $g_d \sim \mathcal{CN}(0, 10^{\frac{PL}{10}})$ ,  $PL = -30.18 - 26 \lg d_s$  and  $d_s$  is the link distance. The channel model from [40] was adopted to model the jamming channel and legitimate channel. The communication scenarios and methods for comparing the simulation are summarized in the table below.

(1) Proposed algorithm. In this scenario, the equilibrium order is chosen as the coalition preference order. The CFG method is used to optimize the clustering, after which the CCD method is used to optimize the phase shift matrix of the RIS.

(2) Selfish order. In this scenario, the selfish order is chosen as the coalition preference order. The deployment method and the optimized algorithm in this scenario are consistent with those of the previous scenario.

(3) Pareto order. In this scenario, the Pareto order is chosen as the coalition preference order. The deployment method and the optimized algorithm in this scenario are consistent with those of the first scenario.

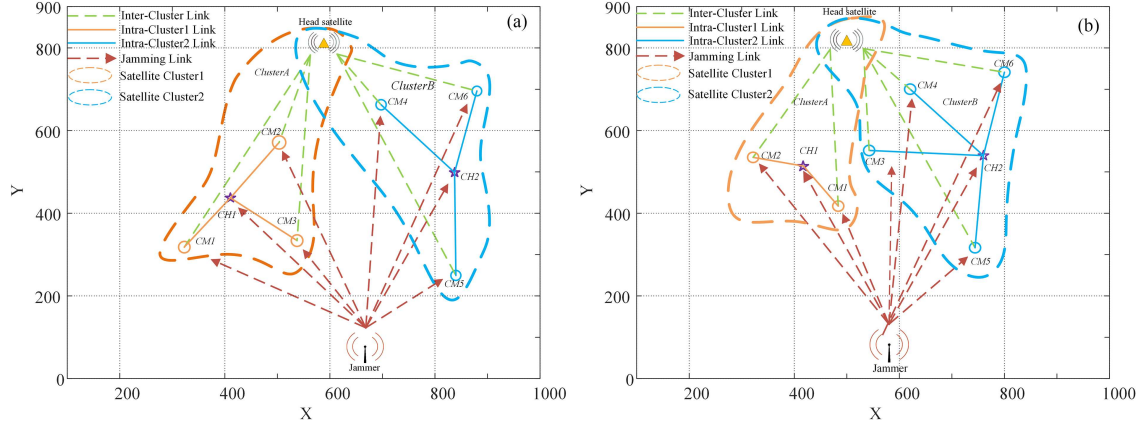
(4) Perfect CSI. In this scenario, the CSI of the imperfect jammer is replaced by the perfect CSI.

(5) SCA algorithm. In this scenario, the coalition preference order is the equilibrium order. The deployment method in this scenario is consistent with that of the previous scenario.

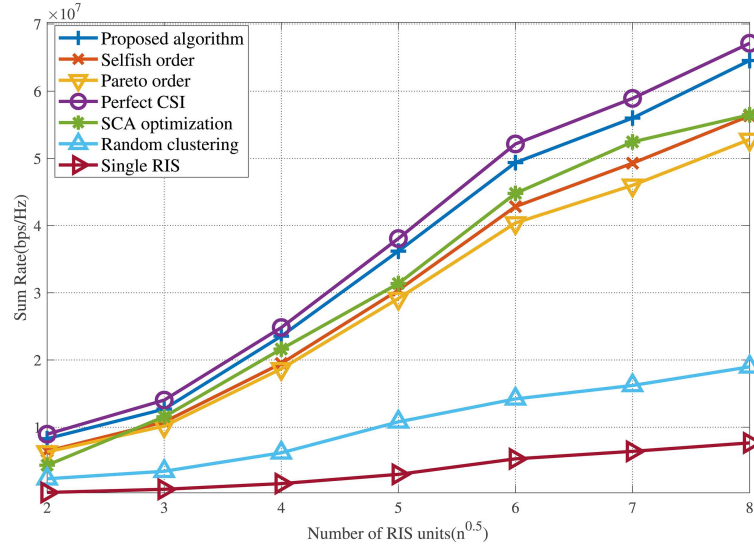
(6) Random clustering. In this scenario, the CFG is no longer used to form clusters. Instead, the RISs are randomly deployed and provided to the shuttle. The other components of the scenario remain unchanged.

(7) Single RIS. In this scenario, only a single intelligent surface, RIS<sub>1</sub>, is retained.

Figure 3 presents the distribution of the satellite RISs and the shuttles using the proposed algorithm in the presence of a jammer. Using the proposed algorithm, the coalition is obtained for each RIS in a distributed way. Figure 3(a) shows the expected condition for satellite clustering, while Figure 3(b)



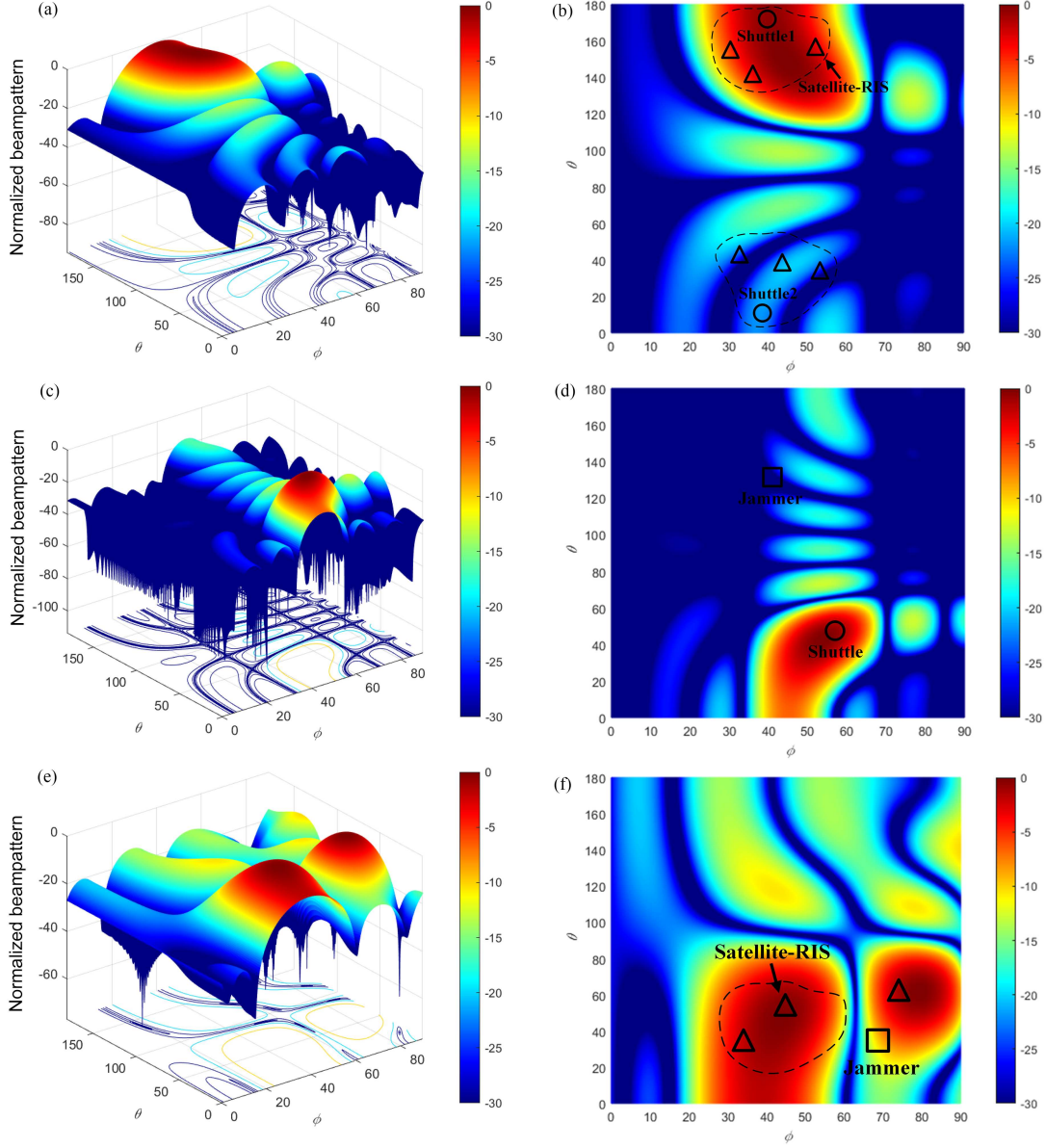
**Figure 3** (Color online) (a) Coalition formation for cohesive clustered satellites under ideal conditions; (b) coalition formation for cohesive clustered satellites with overlapping areas.



**Figure 4** (Color online) Comparison of sum rate of different benchmark schemes and scenarios with different numbers of RIS.

shows the condition where some of the RISs are deployed in overlapping areas. The proposed algorithm could achieve effective clustering under both conditions.

Figure 4 presents a comparison of the performance of the proposed algorithm using different coalition preference orders, beamforming optimization algorithms, and scenarios with different numbers of RIS units. Due to the nonlinear growth of the number of RIS units during the simulation, in order to facilitate analysis and visualization, the value obtained by taking the square root of the number of units is taken as the abscissa in Figure 4. The performance of the proposed algorithm is far superior to that achieved with random clustering. Moreover, the sum rate increases with the number of RIS units because the topology based on game-theoretic clustering efficiently utilizes resources by dynamically adjusting the structure of the satellite. Additionally, the proposed order can achieve better performance compared to the other two orders. The selfish order tends to improve the individual utility when choosing a new coalition, which considers the utility of other coalitions while improving its own utility. In contrast, due to the stringent constraints on the performance of the current coalition and the new coalition, the Pareto order achieves a lower sum rate compared with the equilibrium order. Moreover, for the single RIS with the CCD algorithm, the system performance is significantly lower than that of the proposed algorithm because, in the single RIS scenario, the jammer's channel is positively correlated with the communication channel when there is a jammer. By improving the reflection matrix optimization, the strength of both the received signal and the jammer's signal will be enhanced, resulting in a sharp decline in the communication performance. The simulation result also demonstrates the superior anti-jamming



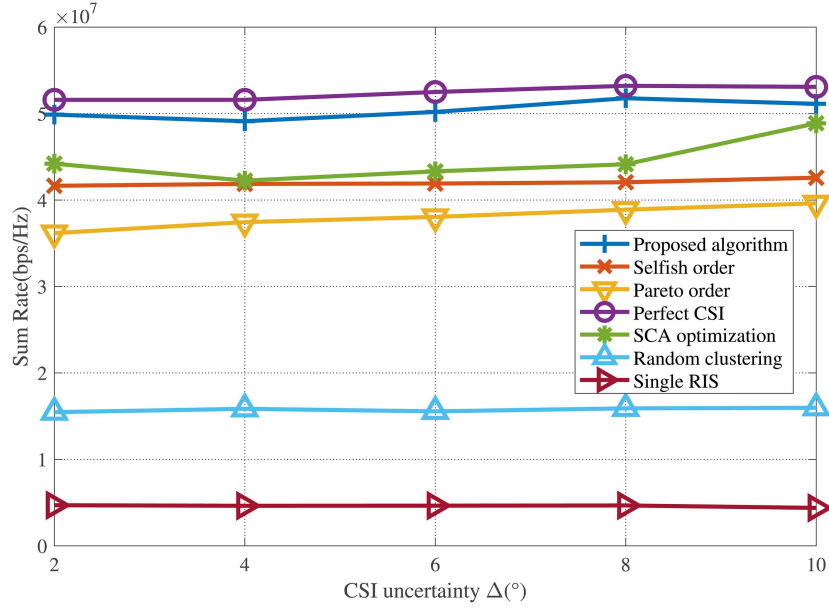
**Figure 5** (Color online) (a) 3D and (b) 2D Tx beam pattern of  $w_{T,1}$ ; (c) 3D and (d) 2D reflecting beam pattern of  $p_2$ ; (e) 3D and (f) 2D Rx beam pattern of  $v_1$ .

performance of the proposed algorithm compared to that of the SCA algorithm.

Figure 5 illustrates the partially optimized radiation beam patterns of the Co1. Here, the main focus is on investigating the Tx beam pattern of HS and the Rx beam pattern of shuttle1, as well as the reflecting beam pattern of the single satellite RIS. The optimized mainlobe of HS clearly points in the region of the target area and simultaneously aligns the nulls toward undesired targets. The beam pattern of the RIS achieved the same performance, the sum rate increased, and the anti-jamming performance was improved. Thus, the proposed game-theoretic clustering and scalable beamforming method can effectively improve the security of transmission and anti-jamming performance under the condition that the jammer's CSI is imperfect.

Figure 6 shows the sum rate of the proposed method for different channel uncertainties, where the data verify the impact of different jamming environments. Here, factors such as the jamming power and location of deployment of the jammer can affect the signal reception of the receiver. Therefore, the comprehensive effect of various jamming factors was considered to verify the performance of the algorithm, namely, the uncertainty of the jamming channel. The performance of all methods remained almost unchanged as the channel uncertainty increased. Robust processing of the imperfect CSI is not





**Figure 6** (Color online) Comparison of the sum rate of different benchmark schemes and scenarios with differences in the CSI uncertainty.

only applicable to the method proposed in this study but also to other methods and can effectively reduce the impact of CSI uncertainty.

## 5 Conclusion

A novel multi-RIS-assisted CCSAJ framework was established to defend against jamming attacks. In a scenario where the jammer's CSI is imperfect, a worst-case optimization problem was formulated to maximize the sum rate by jointly optimizing the satellite clustering, transceiver's beamforming, and phase shifts of the satellite RIS. In particular, the progress of satellite clustering was formulated as a CFG, and the equilibrium order for obtaining the optimal satellite clustering and maximizing the sum rate was proposed. The results prove that the proposed CFG can converge to a stable EPG with at least one NE. An alternative optimization method was also proposed for optimizing the transceiver's beamforming and phase shifts of the satellite RIS. The proposed algorithm has a certain degree of scalability, which is manifested in the following aspects. (1) The CFG can dynamically adjust policies based on the coalition preference order, achieving self-organization of the network. (2) The CCD algorithm reduces the computational complexity, making it applicable to larger-scale scenarios. (3) The modular design enables the system to adapt to constantly changing scenarios. The results of the numerical simulation demonstrate the superiority and validity of the proposed schemes over existing benchmark schemes.

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