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Special Topic: Integration of Large AI Model and  $6\mathrm{G}$ 

# FAS-assisted federated learning over wireless communication systems $^{\dagger}$

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**Abstract** This paper examines the energy efficiency of a multi-user system where federated learning (FL) is implemented in a distributed manner across all nodes. Each user employs a fluid antenna system (FAS) to improve the channel condition, while the base station (BS) is equipped with multiple traditional fixed-position antennas (FPAs). When performing the FL algorithm, each user first trains a local model and transmits it to the BS over shared time-frequency resources. Then, the BS aggregates the received models and broadcasts the combined model back to all users. These steps are repeated until the FL model achieves a desired accuracy level. The system energy is mainly consumed in the computation and transmission processes at the user side. To save energy, we develop an optimization framework that minimizes the total energy consumption by jointly optimizing the learning accuracy, transmit power, antenna positions, and the BS receivers. Since the optimization variables are highly coupled, the problem is non-convex and quite complex. To address the issue, we propose an iterative algorithm to obtain a suboptimal solution of the problem. Simulation results have verified the effectiveness of the algorithm and also the advantages of FAS over the conventional FPA technology.

Keywords fluid antenna system (FAS), federated learning (FL), multiple access, energy efficiency

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# 1 Introduction

The upcoming sixth-generation (6G) wireless networks are expected to provide extremely high capacity, reliability, massive connectivity, and services beyond communications. In line with this trend, fluid antenna system (FAS), representing any software-controllable fluidic, dielectric or conductive structures, such as mechanical liquid-based antennas, radio-frequency pixel-based antennas, movable antennas, or metasurface, that can reconfigure the shape, size, position, orientation, and other radiation characteristics, has been proposed in recent years for enabling flexible and adaptive wireless communications [1–5]. Unlike the traditional antenna techniques where multiple antennas are discretely deployed at fixed positions with sufficient separation, the very fine spatial resolution and dynamic shape of FAS enable it to capitalize on the full range of spatial variations, resulting in significantly improved performance in terms of outage probability [6–8], diversity gain [9,10], system capacity [11–13], and secrecy rate [14,15]. FAS also offers a new capability to exploit the spatial opportunity where the interference suffers from deep fades for multiuser communication, leading to the formation of new multiple access [16–18].

The effectiveness of FAS is closely tied to the accurate acquisition of channel state information (CSI). Therefore, the development of methods for obtaining accurate CSI at different positions or at a wide range of preset ports on a FAS is critical to realizing its full potential. Channel estimation for FAS-assisted systems has been studied by some studies [19,20]. Specifically, in [19], a low-sample size method was proposed to estimate and reconstruct the CSI of a mmWave uplink system, where all transmitters

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use FAS. In [20], channel estimation and reconstruction were achieved through a combination of Nyquist sampling and maximum likelihood estimation. The study highlighted the importance of oversampling to improve estimation accuracy. However, given the practical limitations of oversampling, a suboptimal sampling distance was proposed to enable efficient channel reconstruction while balancing accuracy and feasibility. In addition to the above-mentioned papers, a number of studies have also explored how FAS can work synergistically with other emerging technologies, such as enhancing next-generation multiple access (NGMA) applications [21–23], advancing AI-enabled FAS systems [24–26], investigating FAS-assisted integrated sensing and communication (ISAC) [27–29], and facilitating communication with FAS and reconfigurable intelligent surface (RIS) integration [30,31].

On the other hand, machine learning will be essential in next-generation mobile communication systems, providing the foundation for network intelligence [32–34]. However, due to privacy concerns, limited computational resources, and communication constraints, centralized learning, where the users transmit all of their collected data to a central base station (BS) for processing, is often impractical [35–38]. In response, distributed learning frameworks like federated learning (FL) have emerged as key solutions, allowing users to collaboratively train models while preserving data privacy through local data processing [39, 40]. The performance of FL in wireless systems has been extensively studied in several studies [41–43]. In particular, Ref. [41] addressed the FL loss function minimization problem, factoring in packet errors over wireless links. Studies in [42, 43] focused on minimizing the total energy consumed in computation and communication when executing the FL algorithms. However, Ref. [42] required that all users transmit their local model synchronously. To avoid co-channel interference, Refs. [41, 43] employed orthogonal frequencies for user transmissions, which may constrain the communication performance during model uploads. Moreover, all these studies considered the traditional fixed-position antennas (FPAs).

In this paper, we examine energy-efficient computation and transmission for FL in a FAS-assisted mobile system. Each user has a single-antenna FAS, while the BS is equipped with multiple traditional FPAs. Leveraging collected data, each user trains a local FL model using their own computational resources. After training, users transmit their models to the BS over shared time-frequency resources. The BS aggregates the received models and broadcasts the combined model back to all users. These steps, i.e., local training, data transmission, aggregation, and broadcasting, are repeated iteratively until the FL model achieves a desired accuracy level. The system's energy consumption depends mainly on the learning accuracy and data transmission size. To minimize overall energy usage, we develop an optimization framework that jointly optimizes the learning accuracy, transmit power, antenna positions, and the BS receivers. Due to the strong interdependence among variables, the problem is non-convex and quite intractable. To address this, we develop an iterative algorithm to deal with the problem, where in each iteration, the optimal or near-optimal solutions for the arrived subproblems can be easily found. Simulation results show that the proposed algorithm can greatly reduce the energy consumption, and also highlight the advantages of FAS over the traditional FPA technology.

The rest of this paper is organized as follows. In Section 2, the system model is provided and the problem is formulated. In Section 3, we provide an iterative algorithm to solve the energy minimization problem. Simulation results are given in Section 4 to verify the proposed algorithm. Finally, Section 5 draws some concluding remarks. Auxiliary technical results are given in the appendices.

**Notations.**  $\mathbb{R}$  and  $\mathbb{C}$ , respectively, represent the real and complex spaces. Boldface lower and upper case letters are used to denote vectors and matrices.  $I_M$  stands for the  $M \times M$  dimensional identity matrix and **0** denotes the all-zero vector or matrix. Also, superscripts  $(\cdot)^T$  and  $(\cdot)^H$ , respectively, denote the transpose and conjugate transpose operations. We use calligraphic capital letters to denote sets, such as  $\mathcal{U} = \{1, \ldots, U\}$ , and  $\mathcal{U}_1 \setminus \mathcal{U}_2$  to denote set subtraction. We use calligraphic subscript to denote the set of elements whose indexes take values from the subscript set, e.g.,  $p_{\mathcal{U}} = \{p_1, \ldots, p_U\}$ . If  $|\cdot|$  is applied to a number, it denotes the magnitude of that number; when applied to a set, it represents the cardinality of the set.  $\|\cdot\|_2$  is the Euclidean norm. To help readers follow the mathematical contents, the meanings of some key notations in this paper are summarized in Table 1.

# 2 System model and problem formulation

As depicted in Figure 1, we consider a multi-user wireless communication system with U users and a BS. Each user is equipped with a linear FAS of size  $W_u$ , where the antenna can be instantly switched to

Table 1	List (	of key	notations.
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Notation	Definition	Notation	Definition
$U, U, U_n$	Number of users, set of all users, set of active users at time block $n$	δ	Step size of the gradient method in computing the local FL model
$N, \mathcal{N}, \mathcal{N}_u$	Number of time blocks, set of all time blocks, set of active time blocks of user $u$	$\gamma$	$\gamma$ -strongly convex
$W_u, w_{u,n}$	FAS size of user $u$ , antenna position of the FAS of user $u$ at time block $n$	$\kappa$	Effective switched capacitance in the local computation
M	Number of antennas at the BS	s	Transmit data size of the users
d	Spacing of the BS antenna ULA	$x_{u,n}, p_{u,n}$	Siganl and transmit power of user $u$ at time block $n$
$\lambda$	Wavelength	$\boldsymbol{z}_n, \sigma^2$	Noise vector and noise power
В	Bandwidth	$oldsymbol{y}_n$	Received signal vector of the BS at time block $n$
T	Time constraint on local computation and data transmission	$oldsymbol{g}_{u,n}$	Channel vector from user $u$ to the BS at time block $n$
$ au_n$	Length of time block $n$	$L_u$	Number of propagation paths between user $u$ and the BS
$C_u$	Number of CPU cycles necessary for computing one sample data at user $u$	$\phi_{u,l}$	Complex channel gain of the $l$ -th path of user $u$
$D_u$	Number of data samples of user $u$	$\beta_{u,l}, \theta_{u,l}$	AoA and AoD of the $l$ -th path of user $u$
$G_u$	Computation capacity of user $u$	$a_{u,\mathrm{R}}$	Steering vector at the receiver side
$I_u$	Number of computation iterations of user $u$	$a_{u,\mathrm{T}}$	Steering element at the transmitter side
J	Lipschitz constant for the loss function of the FL model	$oldsymbol{f}_{u,n}$	Linear receiver for detecting $x_{u,n}$
$\eta, \epsilon_0$	Local and global accuracy levels of the FL model	ρ	Fraction of time allocated for local FL model computation



Local FL model

Figure 1 (Color online) FAS-assisted mobile system where the users have FAS for transmission but the BS uses multiple FPAs for reception.

any position in a FAS<sup>1)</sup> [44]. The BS is equipped with an *M*-antenna uniform linear array (ULA) with antenna spacing  $d = \frac{\lambda}{2}$ , where  $\lambda$  is the wavelength. All users and the BS collaboratively execute an FL algorithm across the network for data analysis and inference. The model trained at each user using its local dataset is referred to as the local FL model, while the model aggregated and updated at the BS using all collected local models is called the global FL model. The FL process is conducted iteratively between the users and the BS, with each iteration consisting of three main steps: local computation at

<sup>1)</sup> Each antenna structure can be seen as a reconfigurable pixel-based linear FAS that has many tiny pixels, among which the certain number of pixels and their connections are activated at each time to switch between many reconfigurable states [44]. Using this technology, it is possible to switch the antennas with almost no time delay.

each user, local FL model transmission to the BS, and aggregation of results followed by broadcast from the BS.

Note that in the downlink phase, the BS benefits from high power and bandwidth for broadcasting the global FL model. Thus, as assumed in [43], the downlink transmission time is considered negligible compared to the uplink transmission time. In the following, we first introduce the computation and transmission models, and then formulate the energy-efficiency optimization problem.

#### 2.1 Computation model

In the system, each user u has to compute a local dataset with  $D_u$  data samples. Denote the computation capacity of user u by  $G_u$  (cycles/second), which is measured by the number of CPU cycles per second. Then, the time of user u required for computing its local FL model is

$$t_u^{\rm C} = \frac{I_u C_u D_u}{G_u},\tag{1}$$

where  $C_u$  (cycles/sample) is the number of CPU cycles necessary for computing one sample data at user u. In addition,  $I_u$  in (1) is the number of computation iterations of user u and is given by

$$I_u = v \log_2 \frac{1}{\eta},\tag{2}$$

where  $\eta$  is the local accuracy level and  $v = \frac{2}{(2-J\delta)\delta\gamma}$  [43, Lemma 1]. Here *J* is the Lipschitz constant for the loss function of the FL model,  $\gamma > 0$  is a real positive number that indicates that the loss function of the FL model is  $\gamma$ -strongly convex, and  $\delta < \frac{2}{J}$  is the step size of the gradient method used by each user in computing the local FL model<sup>2</sup> [43]. Then, according to [45, Lemma 1], the energy required to compute a total of  $C_u D_u$  CPU cycles over  $I_u$  iterations is

$$E_u^{\rm C} = \kappa I_u C_u D_u G_u^2, \tag{3}$$

where  $\kappa$  is the effective switched capacitance that depends on the chip architecture.

# 2.2 Transmission model

In the uplink transmission phase, the users communicate with the BS using the same time-frequency resource with bandwidth B. Assume that the users must complete their local computation and data transmission within a total time of T s. As illustrated in Figure 2, user u spends  $t_u^{\rm C}$  on local computation, leaving  $t_u^{\rm T} = T - t_u^{\rm C}$  for data transmission. Since local computations may take different amounts of time for the users, transmissions may thus start at different times for them. Obviously, the user completing its local computation in the shortest time will begin transmission first, while the user with the longest local computation time for the user that completes the computation earliest. This interval is divided into N time blocks  $\tau_1, \ldots, \tau_N$ , with each block marking the start of transmission for a new user who has just finished local computation and begins uploading its FL model. Since there are in total U users, we have at most U transmission time blocks, i.e.,  $N \leq U$ . For convenience, let  $\mathcal{U}_n$  represent the set of users transmitting their FL models during time block n, and let  $\mathcal{N}_u$  denote the set of time block indexes in which user u transmits it FL model.

For ease of understanding, we take the system with three users in Figure 2 for example. Since  $t_2^{\rm C} < t_1^{\rm C} < t_3^{\rm C}$ , the interval  $t_2^{\rm T} = T - t_2^{\rm C}$  is divided into three blocks, i.e.,  $\tau_1 = t_1^{\rm C} - t_2^{\rm C}$ ,  $\tau_2 = t_3^{\rm C} - t_1^{\rm C}$ , and  $\tau_1 = T - t_3^{\rm C}$ . During the first time block  $\tau_1$ , only user 2 is transmitting. Therefore,  $\mathcal{U}_1 = \{2\}$ . Similarly, we have  $\mathcal{U}_2 = \{1, 2\}$  and  $\mathcal{U}_3 = \{1, 2, 3\}$ . On the other hand, user 1 transmits its FL model in the second and third time blocks. Therefore,  $\mathcal{N}_1 = \{2, 3\}$ . Analogously, we have  $\mathcal{N}_2 = \{1, 2, 3\}$  and  $\mathcal{N}_3 = \{3\}$ .

<sup>2)</sup> Note that in this paper, we adopt the same FL algorithm as used in [43]. Therefore, several parameters from [43], such as J,  $\gamma$ , and  $\delta$ , are also relevant here. Due to space limitations, we do not detail the computations for the local and global FL models. For further information on the FL algorithm and the specific meanings of these parameters, please refer to [43].



Figure 2 (Color online) Implementation of the FL algorithm with three users.

#### 2.2.1 Signal model

As stated above, at the *n*-th time block, users in the set  $\mathcal{U}_n$  are transmitting their FL models simultaneously using the shared time-frequency resource. Then, the received signal of the BS at this time block is given by

$$\boldsymbol{y}_n = \sum_{u \in \mathcal{U}_n} \boldsymbol{g}_{u,n} \boldsymbol{x}_{u,n} + \boldsymbol{z}_n, \tag{4}$$

where  $x_{u,n} \sim \mathcal{CN}(0, p_{u,n})$  is the signal of user  $u, p_{u,n}$  is the transmit power,  $g_{u,n} \in \mathbb{C}^{M \times 1}$  denotes the channel vector from user u to the BS, and  $z_n \sim \mathcal{CN}(0, \sigma^2 I_M)$  represents the additive white Gaussian noise (AWGN).

## 2.2.2 Channel model

Let  $w_{u,n} \in [0, W_u]$  denote the antenna position of user u at the *n*-th time block. Using the planar-wave geometric channel model typical in mmWave systems [46], the channel vector  $g_{u,n}$  can be modeled as

$$\boldsymbol{g}_{u,n} = \sqrt{M} \sum_{l=1}^{L_u} \phi_{u,l} \boldsymbol{a}_{u,\mathrm{R}}(\beta_{u,l}) a_{u,\mathrm{T}}^*(\theta_{u,l}, w_{u,n}),$$
(5)

where  $L_u$  is the number of propagation paths between user u and the BS, and  $\phi_{u,l} \sim \mathcal{CN}(0, 10^{-\frac{\mathrm{PL}_{u,l}}{10}})$  is the complex channel gain of the *l*-th path. Here  $\mathrm{PL}_{u,l}$  is the distance-dependent pathloss of the *l*-th path of user u and it is given in the simulation part. Also,  $\boldsymbol{a}_{u,\mathrm{R}}(\beta_{u,l})$  and  $\boldsymbol{a}_{u,\mathrm{T}}(\theta_{u,l}, w_{u,n})$  are, respectively, the steering vector at the receiver side and the steering element at the transmitter side, given by

$$\boldsymbol{a}_{u,\mathrm{R}}(\beta_{u,l}) = \frac{1}{\sqrt{M}} \left[ 1, \mathrm{e}^{-\mathrm{j}\frac{2\pi}{\lambda}d\cos\beta_{u,l}}, \dots, \mathrm{e}^{-\mathrm{j}\frac{2\pi}{\lambda}(M-1)d\cos\beta_{u,l}} \right]^{\mathrm{T}},\tag{6}$$

$$a_{u,\mathrm{T}}(\theta_{u,l}, w_{u,n}) = \mathrm{e}^{-\mathrm{j}\frac{2\pi}{\lambda}w_{u,n}\cos\theta_{u,l}},\tag{7}$$

where  $\beta_{u,l}, \theta_{u,l} \in [0, \pi]$  are, respectively, the angle-of-arrival (AoA) and angle-of-departure (AoD) of the *l*-th path of user *u*. Now, we define

$$\boldsymbol{\Phi}_{u} = \operatorname{diag}\{\phi_{u,1}, \dots, \phi_{u,L_{u}}\} \in \mathbb{C}^{L_{u} \times L_{u}},$$
$$\boldsymbol{A}_{u,\mathrm{R}} = [\boldsymbol{a}_{u,\mathrm{R}}(\beta_{u,1}), \dots, \boldsymbol{a}_{u,\mathrm{R}}(\beta_{u,L_{u}})] \in \mathbb{C}^{M \times L_{u}},$$
$$\boldsymbol{a}_{u,\mathrm{T}}(w_{u,n}) = [a_{u,\mathrm{T}}(\theta_{u,1}, w_{u,n}), \dots, a_{u,\mathrm{T}}(\theta_{u,L_{u}}, w_{u,n})] \in \mathbb{C}^{1 \times L_{u}},$$
(8)

based on which the channel vector  $g_{u,n}$  in (5) can also be expressed as

$$\boldsymbol{g}_{u,n} = \sqrt{M} \boldsymbol{A}_{u,\mathrm{R}} \boldsymbol{\Phi}_{u} \boldsymbol{a}_{u,\mathrm{T}}^{\mathrm{H}}(\boldsymbol{w}_{u,n}).$$
(9)

#### 2.2.3 Wireless transmission

Based on the above settings, the total energy consumed by user u for transmitting its FL model is

$$E_u^{\mathrm{T}} = \sum_{n \in \mathcal{N}_u} \tau_n p_{u,n}.$$
 (10)

Assume that the BS uses a linear receiver  $f_{u,n} \in \mathbb{C}^{M \times 1}$  to detect the symbol  $x_{u,n}$ . Then, according to (4), the data rate of user u at the *n*-th time block is

$$R_{u,n} = \log_2 \left( 1 + \operatorname{SINR}_{u,n} \right),\tag{11}$$

where the signal-to-interference-plus-noise ratio (SINR) is given by

$$\operatorname{SINR}_{u,n} = \frac{p_{u,n} \left| \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{g}_{u,n} \right|^{2}}{\sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} p_{u',n} \left| \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{g}_{u',n} \right|^{2} + \sigma^{2} \left\| \boldsymbol{f}_{u,n} \right\|_{2}^{2}}.$$
(12)

Assuming the dimensions of the users' local FL models are fixed and equal, the data size that each user must upload is thus a constant, denoted by s. Each user u uploads its data over the time blocks in  $\mathcal{N}_u$ . Therefore, the following constraint has to be satisfied:

$$\sum_{n \in \mathcal{N}_u} B\tau_n R_{u,n} \ge s, \ \forall u \in \mathcal{U}.$$
(13)

#### 2.3 Problem formulation

From (3) and (10) we know that the total energy consumption required for local FL model computation and transmission is

$$E = I_0 \sum_{u \in \mathcal{U}} \left( E_u^{\mathrm{C}} + E_u^{\mathrm{T}} \right), \tag{14}$$

where  $I_0 = \frac{b}{1-\eta}$  is the number of global iterations in performing the FL algorithm and  $b = \frac{2J^2}{\gamma^2 \xi} \ln \frac{1}{\epsilon_0}$  [43, Theorem 1]. Here  $\xi$  is a constant value and  $\epsilon_0$  denotes the target accuracy level that the global FL model must achieve. Our objective is to minimize E by optimizing the local accuracy level, users' transmit power, antenna positions, and the BS's linear receivers. The problem can be formulated as

$$\min_{\eta, \boldsymbol{p}_{\mathcal{U}}, \boldsymbol{w}_{\mathcal{U}}, \boldsymbol{F}_{\mathcal{U}}} E \tag{15a}$$

s.t. 
$$t_u^{\mathcal{C}} \leqslant \rho T, \ \forall u \in \mathcal{U},$$
 (15b)

$$0 \leqslant \eta \leqslant 1,\tag{15c}$$

$$\sum_{n \in \mathcal{N}_{u}} B\tau_{n} R_{u,n} \geqslant s, \ \forall u \in \mathcal{U},$$
(15d)

$$p_{u,n} \ge 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u,$$
(15e)

$$0 \leqslant w_{u,n} \leqslant W_u, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u, \tag{15f}$$

$$p_{u,n} = w_{u,n} = 0, \ \boldsymbol{f}_{u,n} = \boldsymbol{0}, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u,$$
 (15g)

where  $0 < \rho < 1$  represents the fraction of time allocated for local FL model computation,  $p_{\mathcal{U}} = \{p_1, \ldots, p_U\}, p_u = [p_{u,1}, \ldots, p_{u,N}]^{\mathrm{T}}, w_{\mathcal{U}} = \{w_1, \ldots, w_U\}, w_u = [w_{u,1}, \ldots, w_{u,N}]^{\mathrm{T}}, F_{\mathcal{U}} = \{F_1, \ldots, F_U\}$ and  $F_u = [f_{u,1}, \ldots, f_{u,N}]$ . Note that in (15g), the transmit power, antenna position, and BS receiver of user u are set to zero for any time block  $n \in \mathcal{N} \setminus \mathcal{N}_u$ . This is because user u only transmits data during the time blocks within  $\mathcal{N}_u$ .

# 3 Iterative algorithm

Since the optimization variables are highly interdependent, problem (15) is quite intractable. In this section, we propose an iterative algorithm to solve it.

# 3.1 Optimization of $\eta$

First, we optimize  $\eta$  with all the other variables fixed. Problem (15) in this case reduces to

$$\min_{\eta} E \tag{16a}$$

s.t. 
$$t_u^{\mathcal{C}} \leqslant \rho T, \ \forall u \in \mathcal{U},$$
 (16b)

$$0 \leqslant \eta \leqslant 1. \tag{16c}$$

To make the analysis more convenient, we first simplify the expression of (16). Let

$$\alpha_1 = \frac{b\kappa v}{\ln 2} \sum_{u \in \mathcal{U}} C_u D_u G_u^2,$$
  

$$\alpha_2 = b \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_u} \tau_n p_{u,n}.$$
(17)

Then, based on (2)–(10), the total energy consumption E in (14) can be rewritten as

$$E = \frac{b}{1-\eta} \sum_{u \in \mathcal{U}} \left( \kappa v C_u D_u G_u^2 \log_2 \frac{1}{\eta} + \sum_{n \in \mathcal{N}_u} \tau_n p_{u,n} \right)$$
$$= \frac{-\alpha_1 \ln \eta + \alpha_2}{1-\eta}.$$
(18)

In addition, according to (1) and (2), the constraint (16b) can be rewritten as

$$-\frac{vC_u D_u}{G_u \ln 2} \ln \eta \leqslant \rho T, \ \forall u \in \mathcal{U},$$
(19)

which can be equivalently transformed to

$$\eta \geqslant \alpha_0,\tag{20}$$

where

$$\alpha_0 = \max\left\{\exp\left(-\frac{\rho T G_u \ln 2}{v C_u D_u}\right), \ \forall u \in \mathcal{U}\right\}.$$
(21)

Since  $\frac{\rho T G_u \ln 2}{v C_u D_u} \ge 0, \forall u \in \mathcal{U}$ , we know that  $0 \le \alpha_0 \le 1$ . Then, problem (16) can be equivalently transformed to

$$\min_{\eta} \quad \frac{-\alpha_1 \ln \eta + \alpha_2}{1 - \eta} \tag{22a}$$

s.t. 
$$\alpha_0 \leqslant \eta \leqslant 1.$$
 (22b)

In Theorem 1, we provide the optimal solution of (22).

**Theorem 1.** Let  $\eta_0$  denote the zero point of  $\frac{1}{\eta} + \ln \eta = 1 + \frac{\alpha_2}{\alpha_1}$ , which is unique and can be easily found using the bisection method. Then, the optimal solution of (22) is

$$\eta^* = \max\{\alpha_0, \eta_0\}. \tag{23}$$

*Proof.* See Appendix A.

# 3.2 Optimization of $p_{\mathcal{U}}, w_{\mathcal{U}}$ , and $F_{\mathcal{U}}$

Next, we optimize  $p_{\mathcal{U}}, w_{\mathcal{U}}$ , and  $F_{\mathcal{U}}$  for given  $\eta$ . The original problem (15) in this case reduces to

$$\min_{\boldsymbol{p}_{\mathcal{U}}, \boldsymbol{w}_{\mathcal{U}}, \boldsymbol{F}_{\mathcal{U}}} \quad \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_{u}} \tau_{n} p_{u, n} \tag{24a}$$

s.t. 
$$\sum_{n \in \mathcal{N}_u} B\tau_n R_{u,n} \ge s, \ \forall u \in \mathcal{U},$$
(24b)

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$$p_{u,n} \ge 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u,$$
(24c)

$$0 \leqslant w_{u,n} \leqslant W_u, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u, \tag{24d}$$

$$p_{u,n} = w_{u,n} = 0, \ \boldsymbol{f}_{u,n} = \boldsymbol{0}, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u,$$
 (24e)

where in the objective function, we omit the constants  $E_u^C$  and  $I_0$  for convenience. From the channel model in (9) and the SINR expression in (12), we know that the optimization variables are coupled, making problem (24) non-convex and quite intractable. To address this issue, we use the fact that if minimum mean square error (MMSE) receivers are adopted by the BS for signal detection, the MMSE and SINR of each link satisfy  $MMSE = \frac{1}{1+SINR}$  [47,48]. This relationship allows us to simplify problem (24). In particular, the mean square error (MSE) of user u at time block n is given by

$$MSE_{u,n} = \mathbb{E}\left[\left|\left(\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{y} - \frac{1}{\sqrt{p_{u,n}}}\boldsymbol{x}_{u,n}\right|^{2}\right]\right]$$
$$= \mathbb{E}\left[\left|\left(\left(\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{g}_{u,n} - \frac{1}{\sqrt{p_{u,n}}}\right)\boldsymbol{x}_{u,n} + \sum_{u'\in\mathcal{U}_{n}\setminus\{u\}}\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{g}_{u',n}\boldsymbol{x}_{u',n} + \boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{z}\right|^{2}\right]\right]$$
$$= \left(\sqrt{p_{u,n}}\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{g}_{u,n} - 1\right)^{2} + \sum_{u'\in\mathcal{U}_{n}\setminus\{u\}}p_{u',n}\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{g}_{u',n}\boldsymbol{g}_{u',n}^{\mathrm{H}}\boldsymbol{f}_{u,n} + \sigma^{2}\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{f}_{u,n}$$
$$= \boldsymbol{f}_{u,n}^{\mathrm{H}}\left(\sum_{u'\in\mathcal{U}_{n}}p_{u',n}\boldsymbol{g}_{u',n}^{\mathrm{H}}\boldsymbol{g}_{u',n} + \sigma^{2}\boldsymbol{I}_{M}\right)\boldsymbol{f}_{u,n} - \sqrt{p_{u,n}}\boldsymbol{f}_{u,n}^{\mathrm{H}}\boldsymbol{g}_{u,n} - \sqrt{p_{u,n}}\boldsymbol{g}_{u,n}^{\mathrm{H}}\boldsymbol{f}_{u,n} + 1, \quad (25)$$

and the corresponding MMSE receiver that minimizes  $MSE_{u,n}$  is

$$\boldsymbol{f}_{u,n}^{*} = \min_{\boldsymbol{f}_{u,n}} \text{MSE}_{u,n}$$
$$= \sqrt{p_{u,n}} \left( \sum_{u' \in \mathcal{U}_n} p_{u',n} \boldsymbol{g}_{u',n} \boldsymbol{g}_{u',n}^{\text{H}} + \sigma^2 \boldsymbol{I}_M \right)^{-1} \boldsymbol{g}_{u,n}.$$
(26)

Since  $MMSE_{u,n} = \frac{1}{1 + SINR_{u,n}}$ , problem (24) can be transformed to

$$\min_{\boldsymbol{p}_{\mathcal{U}}, \boldsymbol{w}_{\mathcal{U}}, \boldsymbol{F}_{\mathcal{U}}} \quad \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_{u}} \tau_{n} p_{u,n} \tag{27a}$$

s.t. 
$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \ln \mathrm{MSE}_{u,n} \leqslant -\hat{s}, \ \forall u \in \mathcal{U},$$
(27b)

$$p_{u,n} \ge 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u,$$
(27c)

$$0 \leqslant w_{u,n} \leqslant W_u, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u, \tag{27d}$$

$$p_{u,n} = w_{u,n} = 0, \ \boldsymbol{f}_{u,n} = \boldsymbol{0}, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u,$$
 (27e)

where  $\hat{s} = \frac{s \ln 2}{B}$ . Note that although the relationship  $\text{MMSE}_{u,n} = \frac{1}{1+\text{SINR}_{u,n}}$  is adopted, we use  $\text{MSE}_{u,n}$  in (27b). Now we explain the reason. From the expressions of the channel vector  $\boldsymbol{g}_{u,n}$  in (9) and the MMSE receiver  $\boldsymbol{f}_{u,n}^*$  in (26), we know that the MMSE receiver minimizing  $\text{MSE}_{u,n}$  is a function of  $p_{u,n}$  and  $w_{u,n}$  for all  $u \in \mathcal{U}_n$ . While substituting  $\boldsymbol{f}_{u,n}^*$  into (25) would allow us to reformulate problem (27) solely in terms of  $\boldsymbol{p}_{\mathcal{U}}$  and  $\boldsymbol{w}_{\mathcal{U}}$ , this approach would result in strong coupling among the variables  $p_{u,n}, w_{u,n}, \forall u \in \mathcal{U}_n$ , rendering problem (27) significantly more intractable. Therefore, we retain  $\boldsymbol{f}_{u,n}$  as an independent variable such that  $\boldsymbol{p}_{\mathcal{U}}, \boldsymbol{w}_{\mathcal{U}}$ , and  $\boldsymbol{F}_{\mathcal{U}}$  can be optimized in an alternative way. Since there is no constraint on the value of  $\boldsymbol{f}_{u,n}$ , we use MSE in (27b). However, it can be easily proven that in the optimal case of (27),  $\boldsymbol{f}_{u,n}$  must adopt the form given in (26), thereby transforming the MSE into the MMSE.

Although the fractional SINR expressions have been avoided, problem (27) remains challenging since the left-hand-side term of each constraint (27b) is a weighted sum of  $|\mathcal{N}_u| \ln(\cdot)$  terms. To remove the  $\ln(\cdot)$  operation, we introduce an auxiliary variable  $\mu_{u,n}$  for each user u and time block n, and arrive at a Xu H, et al. Sci China Inf Sci July 2025, Vol. 68, Iss. 7, 170304:9

new problem as follows:

$$\min_{\boldsymbol{p}_{\mathcal{U}}, \boldsymbol{w}_{\mathcal{U}}, \boldsymbol{F}_{\mathcal{U}}, \boldsymbol{\mu}_{\mathcal{U}}} \quad \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_{u}} \tau_{n} p_{u, n}$$
(28a)

s.t. 
$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \left( e^{\mu_{u,n} - 1} \mathrm{MSE}_{u,n} - \mu_{u,n} \right) \leqslant -\hat{s}, \ \forall u \in \mathcal{U},$$
(28b)

$$p_{u,n} \ge 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u,$$

$$(28c)$$

$$0 \leqslant w_{u,n} \leqslant W_u, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u,$$
(28d)

$$p_{u,n} = w_{u,n} = \mu_{u,n} = 0, \ \boldsymbol{f}_{u,n} = \boldsymbol{0}, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u,$$
(28e)

where  $\boldsymbol{\mu}_{\mathcal{U}} = \{\boldsymbol{\mu}_1, \dots, \boldsymbol{\mu}_U\}$  and  $\boldsymbol{\mu}_u = [\mu_{u,1}, \dots, \mu_{u,N}]^{\mathrm{T}}$ . In Theorem 2, we show that the above two problems (27) and (28) are equivalent and provide the optimal solution for each  $\mu_{u,n}$ .

**Theorem 2.** The problems (27) and (28) are equivalent in that they have the same optimal objective function value. Let  $(\boldsymbol{p}_{\mathcal{U}}^*, \boldsymbol{w}_{\mathcal{U}}^*, \boldsymbol{F}_{\mathcal{U}}^*, \boldsymbol{\mu}_{\mathcal{U}}^*)$  denote the optimal solution of (28) and  $\text{MSE}_{u,n}^*$  represent the corresponding MSE of user u at time block n. Then, we have

$$\mu_{u,n}^* = 1 - \ln \mathrm{MSE}_{u,n}^*, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u.$$

$$\tag{29}$$

#### *Proof.* See Appendix B.

Note that although problem (28) has been significantly simplified, it remains complex. In the following, we solve it by iteratively optimizing the variables.

# 3.2.1 Optimization of $p_{\mathcal{U}}$

We first optimize  $p_{\mathcal{U}}$  for given  $w_{\mathcal{U}}$ ,  $F_{\mathcal{U}}$ , and  $\mu_{\mathcal{U}}$ . For convenience, define  $q_{u,n} = \sqrt{p_{u,n}}, \forall u \in \mathcal{U}, n \in \mathcal{N},$  $q_u = [q_{u,1}, \ldots, q_{u,N}]^{\mathrm{T}}$ , and  $q_{\mathcal{U}} = \{q_1, \ldots, q_U\}$ . Then, based on (25), problem (28) becomes

$$\min_{\boldsymbol{q}_{\mathcal{U}}} \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_{u}} \tau_{n} q_{u,n}^{2} \tag{30a}$$

s.t. 
$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \left\{ e^{\mu_{u,n}-1} \left[ \left( q_{u,n} \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{g}_{u,n} - 1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{2} \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{g}_{u',n} \boldsymbol{g}_{u',n}^{\mathrm{H}} \boldsymbol{f}_{u,n} + \sigma^{2} \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{f}_{u,n} \right] - \mu_{u,n} \right\}$$

$$\leq -\hat{s} \quad \forall u \in \mathcal{U}$$
(30b)

$$q_{u,n} = 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u.$$
(30d)

Problem (30) is a quadratically constrained quadratic program (QCQP) and is convex. The optimal solution can thus be found by using the standard tools such as Interior-point method and CVX. Once the optimal  $q_{\mathcal{U}}^*$  is obtained, we set  $p_{u,n}^* = q_{u,n}^{*2}, \forall u \in \mathcal{U}, n \in \mathcal{N}$ .

# 3.2.2 Optimization of $w_{\mathcal{U}}$ and $F_{\mathcal{U}}$

Next, we optimize  $w_{\mathcal{U}}$  and  $F_{\mathcal{U}}$  with given  $p_{\mathcal{U}}$  and  $\mu_{\mathcal{U}}$ . From (28) we see that the values of  $w_{\mathcal{U}}$  and  $F_{\mathcal{U}}$  do not affect the objective function (28a), but only affect the MSE terms in the constraint (28b). Considering that minimizing the sum MSE of all users at all time blocks can provide more degree of freedom in optimizing the other variables in the subsequent steps, we optimize  $w_{\mathcal{U}}$  and  $F_{\mathcal{U}}$  by minimizing the summation of the left-hand-side terms of (28b) for all users. The problem can be formulated as follows:

$$\min_{\boldsymbol{w}_{\mathcal{U}}, \boldsymbol{F}_{\mathcal{U}}} \quad \sum_{u \in \mathcal{U}} \sum_{n \in \mathcal{N}_u} \tau_n \mathrm{e}^{\mu_{u,n} - 1} \mathrm{MSE}_{u,n} \tag{31a}$$

s.t. 
$$\sum_{n \in \mathcal{N}_u} \tau_n \left( e^{\mu_{u,n} - 1} \mathrm{MSE}_{u,n} - \mu_{u,n} \right) \leqslant -\hat{s}, \ \forall u \in \mathcal{U},$$
(31b)

$$0 \leqslant w_{u,n} \leqslant W_u, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u, \tag{31c}$$

$$w_{u,n} = 0, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N} \setminus \mathcal{N}_u.$$
 (31d)

From the channel model in Subsection 2.2.2 and (25) we observe that the variables  $w_{u,n}$ ,  $f_{u,n}, \forall u \in \mathcal{U}, n \in \mathcal{N}_u$  are highly interdependent in the MSE expressions, making (31) quite challenging to solve. To make the problem tractable, we iteratively optimize the antenna position of each user at each time block, i.e.,  $w_{u,n}$ . Note that although, as seen from (9), the variable  $w_{u,n}$  directly influences only the channel vector  $g_{u,n}$ , the resulting  $g_{u,n}$  subsequently affects the MMSE receivers of all users in the *n*-th time block, i.e.,  $f_{u'',n}, \forall u'' \in \mathcal{U}_n$ . Consequently, in each iteration, in addition to updating  $w_{u,n}$ , it is necessary to also update  $f_{u'',n}, \forall u'' \in \mathcal{U}_n$ . Nevertheless, since  $f_{u'',n}, \forall u'' \in \mathcal{U}_n$  can be computed using the closed-form expression provided in (26), the updating process remains computationally efficient. The problem is formulated as follows:

$$\min_{w_{u,n}, \{\boldsymbol{f}_{u,n}\}_{u'' \in \mathcal{U}_n}} \sum_{u'' \in \mathcal{U}_n} e^{\mu_{u'',n} - 1} MSE_{u'',n}$$
(32a)

t. 
$$\sum_{n'\in\mathcal{N}_{u''}}\tau_{n'}\left(\mathrm{e}^{\mu_{u'',n'}-1}\mathrm{MSE}_{u'',n'}-\mu_{u'',n'}\right)\leqslant -\hat{s}, \ \forall u''\in\mathcal{U}_n,$$
(32b)

$$0 \leqslant w_{u,n} \leqslant W_u, \tag{32c}$$

where the parameter  $\tau_n$  in the objective function is omitted for brevity since it is the weighted sum of the MSE of all users at the fixed time block n. Note that to differentiate from the specific user u and time block n under consideration, in problem (32), we introduce the notation u'' to represent a general user in  $\mathcal{U}_n$  and n' for a general time block in  $\mathcal{N}_{u''}$ . In the following, we show that problem (32) can be near-optimally solved.

Before solving (32), we first simplify its expression. According to (25), for any user  $u'' \in \mathcal{U}_n$ , its MSE at the *n*-th time block can be written as

$$MSE_{u'',n} = \boldsymbol{f}_{u'',n}^{H} \left( \sum_{u' \in \mathcal{U}_{n}} p_{u',n} \boldsymbol{g}_{u',n} \boldsymbol{g}_{u',n}^{H} + \sigma^{2} \boldsymbol{I}_{M} \right) \boldsymbol{f}_{u'',n} - \sqrt{p_{u'',n}} \boldsymbol{f}_{u'',n}^{H} \boldsymbol{g}_{u'',n} - \sqrt{p_{u'',n}} \boldsymbol{g}_{u'',n}^{H} \boldsymbol{f}_{u'',n} \boldsymbol{f}_{u'',n} + 1$$
  
$$= \boldsymbol{f}_{u'',n}^{H} \left( p_{u,n} \boldsymbol{g}_{u,n} \boldsymbol{g}_{u,n}^{H} + \boldsymbol{Q}_{u,n} \right) \boldsymbol{f}_{u'',n} - \sqrt{p_{u'',n}} \boldsymbol{f}_{u'',n}^{H} \boldsymbol{g}_{u'',n} - \sqrt{p_{u'',n}} \boldsymbol{g}_{u'',n}^{H} \boldsymbol{f}_{u'',n} + 1$$
  
$$= -p_{u'',n} \boldsymbol{g}_{u'',n}^{H} \left( p_{u,n} \boldsymbol{g}_{u,n} \boldsymbol{g}_{u,n}^{H} + \boldsymbol{Q}_{u,n} \right)^{-1} \boldsymbol{g}_{u'',n} + 1, \qquad (33)$$

where

$$\boldsymbol{Q}_{u,n} = \sum_{u' \in \mathcal{U}_n \setminus \{u\}} p_{u',n} \boldsymbol{g}_{u',n} \boldsymbol{g}_{u',n}^{\mathrm{H}} + \sigma^2 \boldsymbol{I}_M, \qquad (34)$$

and the last step of (33) holds since in the optimal case,  $f_{u'',n} = \sqrt{p_{u'',n}} \left( p_{u,n} g_{u,n} g_{u,n}^{\mathrm{H}} + Q_{u,n} \right)^{-1} g_{u'',n}$ . Using (33), the objective function (32a) can be rewritten as

$$\sum_{u'' \in \mathcal{U}_n} e^{\mu_{u'',n} - 1} MSE_{u'',n} = \sum_{u'' \in \mathcal{U}_n} e^{\mu_{u'',n} - 1} \left[ -p_{u'',n} \boldsymbol{g}_{u'',n}^{\mathrm{H}} \left( p_{u,n} \boldsymbol{g}_{u,n} \boldsymbol{g}_{u,n}^{\mathrm{H}} + \boldsymbol{Q}_{u,n} \right)^{-1} \boldsymbol{g}_{u'',n} + 1 \right].$$
(35)

In addition, the constraint (32b) can be rewritten as

 $\mathbf{S}$ 

$$\sum_{n' \in \mathcal{N}_{u''}} \tau_{n'} \left( e^{\mu_{u'',n'} - 1} \mathrm{MSE}_{u'',n'} - \mu_{u'',n'} \right)$$

$$= \tau_{n} \left( e^{\mu_{u'',n} - 1} \mathrm{MSE}_{u'',n} - \mu_{u'',n} \right) + \sum_{n' \in \mathcal{N}_{u''} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u'',n'} - 1} \mathrm{MSE}_{u'',n'} - \mu_{u'',n'} \right)$$

$$= \tau_{n} \left\{ e^{\mu_{u'',n} - 1} \left[ -p_{u'',n} g_{u'',n}^{\mathrm{H}} \left( p_{u,n} g_{u,n} g_{u,n}^{\mathrm{H}} + Q_{u,n} \right)^{-1} g_{u'',n} + 1 \right] - \mu_{u'',n} \right\}$$

$$+ \sum_{n' \in \mathcal{N}_{u''} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u'',n'} - 1} \mathrm{MSE}_{u'',n'} - \mu_{u'',n'} \right)$$

$$\leqslant -\hat{s}, \ \forall u'' \in \mathcal{U}_{n}. \tag{36}$$

For convenience, we equivalently rewrite the inequality (36) as follows:

$$-p_{u^{\prime\prime},n}\boldsymbol{g}_{u^{\prime\prime},n}^{\mathrm{H}}\left(p_{u,n}\boldsymbol{g}_{u,n}\boldsymbol{g}_{u,n}^{\mathrm{H}}+\boldsymbol{Q}_{u,n}\right)^{-1}\boldsymbol{g}_{u^{\prime\prime},n}\leqslant\psi_{u^{\prime\prime},n},\;\forall u^{\prime\prime}\in\mathcal{U}_{n},$$
(37)

Algorithm 1 Iterative optimization for solving (31).

1: repeat

2: **for** n = 1 : N **do** 

- 3: for  $u \in \mathcal{U}_n$  do
- 4: Find the point in  $\{0, \varepsilon, 2\varepsilon, \ldots, W_u\}$  that minimizes (39a) and satisfies (39b) and (39c) using one-dimensional search, and let it be  $w_{u,n}^*$ ;
- 5: Update  $g_{u,n}$  based on  $w_{u,n}^*$  and (9);
- 6: Update  $f_{u,n}, \forall u'' \in \mathcal{U}_n$  based on (26) and the new  $g_{u,n}$ ;
- 7: end for
- 8: end for
- 9: **until** Convergence;

#### Algorithm 2 Iterative optimization for solving (15).

- 1: Initialize  $p_{\mathcal{U}}$  and  $w_{\mathcal{U}}$ ;
- 2: Compute  $F_{\mathcal{U}}$  and  $\mu_{\mathcal{U}}$  based on (26) and Theorem 2;
- 3: repeat
- 4: Update  $\eta$  based on Theorem 1;
- 5: repeat
- 6: Update  $q_{\mathcal{U}}$  by solving (30) and set  $p_{u,n} = q_{u,n}^2, \forall u \in \mathcal{U}, n \in \mathcal{N};$
- 7: Update  $w_{\mathcal{U}}$  and  $F_{\mathcal{U}}$  using Algorithm 1;
- 8: Update  $\mu_{\mathcal{U}}$  based on Theorem 2;
- 9: **until** Convergence;
- 10: **until** Convergence;

where

$$\psi_{u'',n} = \frac{1}{\tau_n \mathrm{e}^{\mu_{u'',n}-1}} \left[ \tau_n \mu_{u'',n} - \sum_{n' \in \mathcal{N}_{u''} \setminus \{n\}} \tau_{n'} \left( \mathrm{e}^{\mu_{u'',n'}-1} \mathrm{MSE}_{u'',n'} - \mu_{u'',n'} \right) - \hat{s} \right] - 1.$$
(38)

Based on (35) and (37), the problem (32) can be equivalently transformed to

$$\max_{w_{u,n}} \sum_{u'' \in \mathcal{U}_n} e^{\mu_{u'',n} - 1} p_{u'',n} \boldsymbol{g}_{u'',n}^{\mathrm{H}} \left( p_{u,n} \boldsymbol{g}_{u,n} \boldsymbol{g}_{u,n}^{\mathrm{H}} + \boldsymbol{Q}_{u,n} \right)^{-1} \boldsymbol{g}_{u'',n}$$
(39a)

s.t. 
$$-p_{u'',n}\boldsymbol{g}_{u'',n}^{\mathrm{H}}\left(p_{u,n}\boldsymbol{g}_{u,n}\boldsymbol{g}_{u,n}^{\mathrm{H}}+\boldsymbol{Q}_{u,n}\right)^{-1}\boldsymbol{g}_{u'',n} \leqslant \psi_{u'',n}, \ \forall u'' \in \mathcal{U}_n,$$
 (39b)

$$0 \leqslant w_{u,n} \leqslant W_u. \tag{39c}$$

We divide the interval  $[0, W_u]$  into K + 1 equal-length subintervals, yielding K + 1 boundary points  $\{0, \varepsilon, 2\varepsilon, \ldots, W_u\}$ , where  $\varepsilon = W_u/K$  is the length of each subinterval. Then, a near-optimal solution of (39) can be easily found by using the exhaustive search (ES) (or one-dimensional search) method<sup>3)</sup>. Once the near-optimal  $w_{u,n}^*$  is obtain, we update  $g_{u,n}$  using (9), and then  $f_{u'',n}, \forall u'' \in \mathcal{U}_n$  based on (26). The iterative steps for solving (31) are summarized in Algorithm 1.

# 3.2.3 Optimization of $\mu_{\mathcal{U}}$

Given  $p_{\mathcal{U}}$ ,  $w_{\mathcal{U}}$ , and  $F_{\mathcal{U}}$ , the optimal  $\mu_{\mathcal{U}}^*$  can be obtained based on Theorem 2. So far, we have shown how to optimize the variables in each iteration. The main steps for solving the original problem (15) are summarized in Algorithm 2.

# 3.3 Convergence and complexity analysis

In this subsection, we analyze the convergence and complexity of the proposed Algorithms 1 and 2.

#### 3.3.1 Convergence analysis

As shown in Algorithm 1, for each time block n and user u in  $\mathcal{U}_n$ , a near-optimal  $w_{u,n}$  can be obtained by solving (39). Then,  $f_{u,n}, \forall u'' \in \mathcal{U}_n$  are optimally updated based on (26). Therefore, in each iteration, the objective function of (31) decreases. Since the MSE is obviously limited, the convergence of Algorithm 1 is guaranteed. In Algorithm 2,  $\eta$ ,  $q_{\mathcal{U}}$ ,  $(w_{\mathcal{U}}, F_{\mathcal{U}})$ , and  $\mu_{\mathcal{U}}$  are iteratively updated. In each step, the

<sup>3)</sup> Here we say that the solution is near-optimal since the variable  $w_{u,n}$  is continuous and it is impossible to really perform ES. Therefore, we discretize the interval and carry out the one-dimensional search, whose performance depends on the size of the search step  $\varepsilon$ .

optimal or near-optimal solutions of these variables can be obtained, making the corresponding objective function decrease. Since the total energy consumption is limited, Algorithm 2 thus converges.

#### 3.3.2 Complexity analysis

To assess the computational complexity of the proposed algorithms, we quantify the total number of floating-point operations (FLOPs), where one FLOP corresponds to either a complex multiplication or a complex summation. The complexity is expressed as a polynomial function of the dimensions of the involved matrices. For simplicity, we retain only the leading (i.e., highest-order or dominant) terms in the polynomial, as these terms primarily determine the asymptotic behavior of the complexity.

The complexity of Algorithm 1 primarily arises from the channel inversions performed in step 4, which involves a one-dimensional search to identify the near-optimal antenna position  $w_{u,n}$ . In this step, the objective function and constraints of (39) have to be computed K + 1 times, with each computation requiring  $|\mathcal{U}_n|$  channel inversions. Given that the inversion of an  $M \times M$  Hermitian matrix requires  $\mathcal{O}(M^3)$  FLOPs, each inner iteration of Algorithm 1 incurs a complexity of  $\mathcal{O}(M^3|\mathcal{U}_n|(K+1))$ . Algorithm 1 consists of N outer iterations, with each outer iteration involving  $|\mathcal{U}_n|$  inner iterations. Since  $|\mathcal{U}_n|$  varies with n and satisfies  $|\mathcal{U}_n| \leq |\mathcal{U}| = U$ , we simplify the complexity analysis by replacing  $|\mathcal{U}_n|$  with U. Then, the overall complexity of Algorithm 1 is  $\mathcal{O}(M^3U^2N(K+1))$ . In Algorithm 2, the main complexity lies in step 7, i.e., using Algorithm 1 to update  $w_{\mathcal{U}}$  and  $F_{\mathcal{U}}$ . Let  $\nu_1$  and  $\nu_2$  denote the numbers of outer and inner iterations of Algorithm 2. Then, Algorithm 2 involves a total complexity of  $\mathcal{O}(\nu_1\nu_2M^3U^2N(K+1))$ .

# 4 Simulation results

In this section, we evaluate the performance of the considered system and also the proposed algorithms by simulation. We consider the UMa (Urban Macro) environment. According to [49], the distance-dependent pathloss of the *l*-th path of user *u* is  $PL_{u,l} = 32.4 + 20 \log_{10} f_c + 30 \log_{10} \Delta_u + \varsigma_{u,l}$ , where  $\Delta_u$  is the distance in meters between user *u* and the BS, and  $\varsigma_{u,l} \sim \mathcal{N}(0, \sigma_{SF}^2)$  is the log-normal shadow fading.  $\sigma_{SF}$  is the standard deviation of the shadow fading and is set to be 7.8 dB here. The mmWave system operates at a carrier frequency  $f_c = 28$  GHz and the bandwidth is B = 20 MHz. The noise power is  $\sigma^2 = -94$  dBm. In addition, we set J = 10,  $\delta = 1/10$ ,  $\gamma = 1$ ,  $\kappa = 10^{-29}$ ,  $\xi = 1/10$ ,  $\epsilon_0 = 10^{-3}$ , and  $\rho = 1/2$ . The parameter  $C_u$  is uniformly distributed in  $[10^4, 3 \times 10^4]$  cycles/sample. Each user has  $D_u = 500$  data samples, which are randomly selected from the dataset with equal probability. The computation capacity of each user is  $G_u = 1$  GHz. For convenience, we assume equal FAS size and number of paths for all users, i.e.,  $W_u = W$ ,  $L_u = L, \forall u \in \mathcal{U}$ . The other parameters are specified in each figure. We compare the performance of FAS with the FPA strategy, which follows the same steps as Algorithm 2, but omits updates to  $w_{\mathcal{U}}$ . All results are obtained by averaging over 1000 independent channel realizations, with each realization generated randomly according to the specified distribution.

Figure 3 depicts the convergence behaviors of the proposed Algorithm 2 under different configurations of U. It can be seen from this figure that the total energy consumption decreases monotonically during the iterative procedure and converges rapidly in about 4 iterations for all configurations. Moreover, Figure 3 shows that the energy consumption grows with U. This is consistent with intuition since more users result in more computation load, transmission data, and co-channel interference.

Figure 4 investigates the effect of the number of BS antennas and the transmit data size of the users. Several observations can be made from this figure. First, the total energy consumption decreases with M. This is because when the BS has more antennas, it becomes more powerful in receiving the desired signal and compressing the co-channel interference. Second, it can be seen that when M is small, in contrast to the FPA scheme, FAS can dramatically decrease the energy consumption. For example, when M = 2 and s = 60 kbits. FPA requires about 660 J to complete the task, while FAS requires only 420 J. However, as the number of antennas increases, this advantage becomes less pronounced. This suggests that FAS and the multi-antenna techniques can be complementary. In addition, we see that as expected, the energy consumption increases with the transmit data size s, and when FAS demonstrates a more significant advantage over FPA in terms of the energy efficiency.

Figure 5 compares the performance of FAS with that of FPA under different configurations of L and U. Besides the observation that the energy consumption increases with U, which has already been verified by Figure 3, Figure 5 also shows that the energy consumption decreases with L, indicating that a rich scattering environment enhances communication effectiveness. In addition, it can be seen that the gap



Figure 3 (Color online) Convergence behaviors of Algorithm 2 with  $W = 10\lambda$ , M = 4, L = 5, and s = 60 kbits.



Figure 5 (Color online) Total energy consumption versus the number of paths with  $W = 10\lambda$ , M = 4, and s = 60 kbits.



**Figure 4** (Color online) Total energy consumption versus the number of BS antennas with U = 4,  $W = 10\lambda$ , and L = 5.



Figure 6 (Color online) Total energy consumption versus the normalized size of FAS with U = 4, M = 4, and L = 5.

between FAS and FPA first increases and then decreases with L. Interestingly, when L = 1, i.e., there is only one path between each user and the BS, the energy consumption required by the methods using FAS and FPA is exactly the same. Therefore, using FAS in this case does not benefit the performance. Note that similar observations have also been made in the sum capacity maximization problem of a multi-user uplink communication system [11], and have been theoretically analyzed (see [11, Lemma 1]).

Figure 6 investigates the effect of the FAS size W. Note that W does not affect the performance of the FPA method. Therefore, for any given s, the total energy consumed by the FPA method is a constant. It is seen from this figure that as W increases, the total energy consumed by the FAS strategy decreases significantly at the beginning and then saturates. This is because we use the same K = 100for all configurations when solving the problem (39) in Algorithm 1. Then, when W is relatively small, increasing W can significantly reduce the correlation among the K + 1 sampling points. However, when W becomes large, the spatial correlation among the K + 1 sampling points is no longer significant. In this case, for a fixed K, further increases in W will not result in a significant reduction in energy consumption.

# 5 Conclusion

This paper investigated FAS-assisted FL in a multi-user communication system, aiming to improve the energy efficiency of the FL process. We addressed this by formulating an optimization problem that minimizes the total energy consumed for user-side computation and data transmission, subject to latency and data size constraints. Due to the coupling of variables such as learning accuracy, transmit power, antenna positions, and BS receivers, the considered problem is inherently non-convex and challenging to

solve. To overcome this, we proposed a low-complexity algorithm that works iteratively and provides a near-optimal solution. Simulation results have confirmed the algorithm's effectiveness and highlighted the advantages of FAS over FPA.

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#### Appendix A Proof of Theorem 1

According to (18), the first-order derivative of E w.r.t.  $\eta$  is

$$\frac{\partial E}{\partial \eta} = \frac{-\alpha_1(\frac{1}{\eta} + \ln \eta) + \alpha_1 + \alpha_2}{(1-\eta)^2}.$$
(A1)

For convenience, denote function  $r(\eta) = \frac{1}{\eta} + \ln \eta$ . Its first-order derivative is given by

$$\frac{\partial r}{\partial \eta} = \frac{\eta - 1}{\eta^2},\tag{A2}$$

which is non-positive in (0, 1] and non-negative in  $[1, +\infty)$ . Therefore, when  $\eta$  increases from 0 to 1,  $r(\eta)$  decreases monotonically from  $+\infty$  to 1. Since  $\alpha_1$  and  $\alpha_2$  are both positive, the numerator of (A1) increases monotonically from  $-\infty$ to  $\alpha_2$  as  $\eta$  increases from 0 to 1. Then, there must exist a unique zero point of the numerator of (A1) in the interval (0, 1]. Let  $\eta_0 \in (0, 1]$  denote this zero point, i.e., it makes the following equation holds:

$$\frac{1}{\eta} + \ln \eta = 1 + \frac{\alpha_2}{\alpha_1}.\tag{A3}$$

Since  $\frac{1}{n} + \ln \eta$  decreases monotonically w.r.t.  $\eta$  in (0, 1],  $\eta_0$  can be easily found using the bisection method.

Accordingly, we know that  $\partial E/\partial \eta$  is non-positive in  $(0, \eta_0]$  and non-negative in  $[\eta_0, +\infty)$ , indicating that E decreases w.r.t.  $\eta$  in  $(0, \eta_0]$  and then increases in  $[\eta_0, +\infty)$ . Since  $\eta_0 \in (0, 1]$  and  $\eta$  belongs to  $[\alpha_0, 1]$  (see constraint (22b)), we know that if  $\alpha_0 < \eta_0$ , the optimal solution of (22) is  $\eta_0$ . Otherwise, it is  $\alpha_0$ . Theorem 1 is thus proven.

## Appendix B Proof of Theorem 2

Let  $(\boldsymbol{p}_{\mathcal{U}}^{*}, \boldsymbol{w}_{\mathcal{U}}^{*}, \boldsymbol{F}_{\mathcal{U}}^{*})$  denote the optimal solution of (27) and  $\text{MSE}_{u,n}^{*}$  represent the corresponding MSE of user u at time block n. Furthermore, let  $\mu_{u,n}^{*} = 1 - \ln \text{MSE}_{u,n}^{*}$ . Substituting  $\mu_{u,n}^{*}$  and  $\text{MSE}_{u,n}^{*}$  into (28b), we have

$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \left( e^{\mu_{u,n}^{*} - 1} \mathrm{MSE}_{u,n}^{*} - \mu_{u,n}^{*} \right) = \sum_{n \in \mathcal{N}_{u}} \tau_{n} \ln \mathrm{MSE}_{u,n}^{*}$$
$$\leqslant -\hat{s}, \ \forall u \in \mathcal{U}, \tag{B1}$$

where the last step follows from the fact that  $MSE_{u,n}^*$  makes (27b) hold. Then, we know that  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  can make all the constraints of (28) satisfied and is thus a solution of (28).

Let  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  denote the optimal solution of (28) and  $MSE_{u,n}^*$  represent the corresponding MSE of user u at time block n. In the following we first show that (29) is true, based on which we then know that  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*)$  is a solution of (27). Based on (25), the constraint (28b) can be rewritten as

$$\sum_{n \in \mathcal{N}_u} \tau_n \left( e^{\mu_{u,n} - 1} \mathrm{MSE}_{u,n} - \mu_{u,n} \right)$$

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$$= \sum_{n \in \mathcal{N}_{u}} \tau_{n} \left\{ e^{\mu_{u,n}^{*} - 1} \left[ \left( \sqrt{p_{u,n}^{*}} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*} - 1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} p_{u',n}^{*} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{*} \right\} \\ \leqslant -\hat{s}, \ \forall u \in \mathcal{U}.$$
(B2)

For convenience, we define  $q_{u,n}^* = \sqrt{p_{u,n}^*}$  and rewrite (B2) equivalently as follows:

$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \left\{ e^{\mu_{u,n}^{*} - 1} \left[ \left( q_{u,n}^{*} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*} - 1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{*2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{*} \right\} \\ \leqslant -\hat{s}, \ \forall u \in \mathcal{U}.$$
(B3)

To prove (29), we first give Lemma B1, which provides an upper bound on the optimal transmit power. Lemma B1. In the optimal case of problem (28), the transmit power satisfies

$$\sqrt{p_{u,n}^*} = q_{u,n}^* \leqslant \frac{1}{\boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^*}, \ \forall u \in \mathcal{U}, \ n \in \mathcal{N}_u.$$
(B4)

Proof. See Appendix C.

Now we prove (29) by contraposition. Assume that Eq. (29) is not true. Then, there must exist a user  $u \in \mathcal{U}$  and a time block  $n \in \mathcal{N}_u$  such that

$$\mu_{u,n}^* \neq 1 - \ln \mathrm{MSE}_{u,n}^*. \tag{B5}$$

For this given user  $u \in \mathcal{U}$ , the corresponding inequality in (B3) can be rewritten as

$$\tau_{n} \left\{ e^{\mu_{u,n}^{*}-1} \left[ \left( q_{u,n}^{*} \boldsymbol{f}_{u,n}^{*H} \boldsymbol{g}_{u,n}^{*}-1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{*2} \boldsymbol{f}_{u,n}^{*H} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*H} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*H} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{*} \right\} + \sum_{n' \in \mathcal{N}_{u} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u,n'}^{*}-1} \mathrm{MSE}_{u,n'}^{*} - \mu_{u,n'}^{*} \right) \leqslant -\hat{s},$$
(B6)

in which we separate the summation terms at time block n from those at time blocks  $n' \in \mathcal{N}_u \setminus \{n\}$  for convenience. Note that since  $MSE_{u,n}^* \ge 0$ ,  $e^{\mu_{u,n}-1}MSE_{u,n}^* - \mu_{u,n}$  is a convex function of  $\mu_{u,n}$  and its unique optimal solution is

$$u_{u,n}^{\star} = \arg \min_{\mu_{u,n}} \left( e^{\mu_{u,n} - 1} MSE_{u,n}^{*} - \mu_{u,n} \right)$$
  
= 1 - ln MSE<sub>u,n</sub><sup>\*</sup>. (B7)

Then, by replacing  $\mu_{u,n}^*$  in (B6) with  $\mu_{u,n}^*$ , we have

$$\tau_{n} \left\{ e^{\mu_{u,n}^{*}-1} \left[ \left( q_{u,n}^{*} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*}-1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{*2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{*} \right\} + \sum_{n' \in \mathcal{N}_{u} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u,n'}^{*}-1} \mathrm{MSE}_{u,n'}^{*} - \mu_{u,n'}^{*} \right) < -\hat{s}.$$
(B8)

Note that the inequality in (B8) strictly holds, while in (B6), the inequality may hold with equality. This is because  $\mu_{u,n}^*$  is the unique point that minimizes  $e^{\mu_{u,n}-1} MSE_{u,n}^* - \mu_{u,n}$  (see (B7)), while  $\mu_{u,n}^*$  cannot due to (B5). As discussed in Appendix C, the term  $\tau_n e^{\mu_{u,n}^* - 1} \left( q_{u,n}^* f_{u,n}^{*H} g_{u,n}^* - 1 \right)^2$  in (B8) is a quadratic function of  $q_{u,n}^*$  with an upward opening and its zero point is  $\frac{1}{f_{u,n}^{*H} g_{u,n}^*}$ . Since  $q_{u,n}^* \leqslant \frac{1}{f_{u,n}^{*H} g_{u,n}^*}$  (see Lemma B1), we can always find a new  $q_{u,n}^*$  such that

$$q_{u,n}^{\star} < q_{u,n}^{*},\tag{B9}$$

and

$$\tau_{n} \left\{ e^{\mu_{u,n}^{\star} - 1} \left[ \left( q_{u,n}^{\star} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*} - 1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{*2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{\star} \right\} + \sum_{n' \in \mathcal{N}_{u} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u,n'}^{*} - 1} \mathrm{MSE}_{u,n'}^{*} - \mu_{u,n'}^{*} \right) = -\hat{s}.$$
(B10)

Note that as discussed in Appendix C, for given  $u \in \mathcal{U}$  and  $n \in \mathcal{N}_u$ , besides the inequality in (B3) with u, which is exactly (B8),  $q_{u,n}^*$  also exists in  $|\mathcal{U}_n| - 1$  other inequalities in (B3) with  $u'' \in \mathcal{U}_n \setminus \{u\}$  (see (C2)). In particular, for each  $u'' \in \mathcal{U}_n \setminus \{u\}$ ,  $q_{u,n}^*$  exists in the term  $\tau_n e^{\mu_u^* \prime \prime}, n^{-1} q_{u,n}^{*2} f_{u'',n}^{*H} g_{u,n}^* g_{u,n}^{*H} f_{u'',n}^*$ . Due to (B9), it is obvious that

$$\tau_{n} e^{\mu_{u'',n}^{*}-1} q_{u,n}^{*2} \boldsymbol{f}_{u'',n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*} \boldsymbol{g}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u'',n}^{*} < \tau_{n} e^{\mu_{u'',n}^{*}-1} q_{u,n}^{*2} \boldsymbol{f}_{u'',n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*} \boldsymbol{g}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u'',n}^{*}, \quad \forall u'' \in \mathcal{U}_{n} \setminus \{u\}.$$
(B11)

Therefore, we know that by replacing  $p_{u,n}^*$  and  $\mu_{u,n}^*$  in  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  with  $p_{u,n}^* = q_{u,n}^{*2}$  and  $\mu_{u,n}^*$ , we can obtain a new feasible solution for (28), and due to (B9), the objective function value can be further decreased. This contradicts the assumption that  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  is the optimal solution of (28). Therefore, Eq. (29) is true.

On the other hand, due to (29), we have

$$\sum_{n \in \mathcal{N}_{u}} \tau_{n} \left( e^{\mu_{u,n}^{*} - 1} \mathrm{MSE}_{u,n}^{*} - \mu_{u,n}^{*} \right) = \sum_{n \in \mathcal{N}_{u}} \tau_{n} \ln \mathrm{MSE}_{u,n}^{*}$$
$$\leq -\hat{s}, \ \forall u \in \mathcal{U}.$$
(B12)

Therefore,  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*)$  is also a solution of (28). Obviously, the optimal objective function value of (27) resulted from  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*)$  is equal to that of (28) resulted from  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$ , since otherwise the objective function value can be further decreased by obtaining a new solution for (27) from  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  or obtaining a new solution for (28) from  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  or obtaining a new solution for (28) from  $(p_{\mathcal{U}}^*, w_{\mathcal{U}}^*, F_{\mathcal{U}}^*, \mu_{\mathcal{U}}^*)$  or obtaining a new solution. Theorem 2 is thus proven.

## Appendix C Proof of Lemma B1

For any given time block  $n \in \mathcal{N}$  and user  $u \in \mathcal{U}_n$ , from (B3) we know that only inequalities with  $u \in \mathcal{U}_n$  contain  $q_{u,n}^*$ . This is because, during the time block n, only users in  $\mathcal{U}_n$  are transmitting their signals. Therefore, the signal of user u only affects the communication of users in  $\mathcal{U}_n$ . Then,  $q_{u,n}^*$  appears in  $|\mathcal{U}_n|$  inequalities in (B3), i.e., the  $\tau_n e^{\mu_{u,n}^* - 1} \left( q_{u,n}^* f_{u,n}^{*H} g_{u,n}^* - 1 \right)^2$  term in

$$\tau_{n} e^{\mu_{u,n}^{*}-1} \left(q_{u,n}^{*} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u,n}^{*}-1\right)^{2} + \tau_{n} \left\{ e^{\mu_{u,n}^{*}-1} \left[ \sum_{u' \in \mathcal{U}_{n} \setminus \{u\}} q_{u',n}^{*2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{g}_{u',n}^{*} \boldsymbol{g}_{u',n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} + \sigma^{2} \boldsymbol{f}_{u,n}^{*\mathrm{H}} \boldsymbol{f}_{u,n}^{*} \right] - \mu_{u,n}^{*} \right\} + \sum_{n' \in \mathcal{N}_{u} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u,n'}^{*}-1} \mathrm{MSE}_{u,n'}^{*} - \mu_{u,n'}^{*} \right) \leqslant -\hat{s},$$
(C1)

and the  $\tau_n e^{\mu_{u'',n}^* - 1} q_{u,n}^{*2} f_{u'',n}^{*H} g_{u,n}^* g_{u,n}^{*H} f_{u'',n}^*$  term in

$$\tau_{n} e^{\mu_{u'',n}^{*} - 1} q_{u,n}^{*2} f_{u'',n}^{*H} g_{u,n}^{*} g_{u,n}^{*H} f_{u'',n}^{*} + \tau_{n} \left\{ e^{\mu_{u'',n}^{*} - 1} \left[ \left( q_{u'',n}^{*} f_{u'',n}^{*H} g_{u'',n}^{*} - 1 \right)^{2} + \sum_{u' \in \mathcal{U}_{n} \setminus \{u'',u\}} q_{u',n}^{*2} f_{u'',n}^{*H} g_{u',n}^{*} g_{u',n}^{*H} f_{u'',n}^{*} + \sigma^{2} f_{u'',n}^{*H} f_{u'',n}^{*} \right] - \mu_{u'',n}^{*} \right\} + \sum_{n' \in \mathcal{N}_{u''} \setminus \{n\}} \tau_{n'} \left( e^{\mu_{u'',n'}^{*} - 1} \mathrm{MSE}_{u'',n'}^{*} - \mu_{u'',n'}^{*} \right) \leqslant -\hat{s}, \ \forall u'' \in \mathcal{U}_{n} \setminus \{u\}.$$
(C2)

From (26) we know that  $f_{u,n}^{*H}g_{u,n}^{*}$  is real and positive. The zero point of

$$\tau_n \mathrm{e}^{\mu_{u,n}^* - 1} \left( q_{u,n}^* \boldsymbol{f}_{u,n}^{\mathrm{H}} \boldsymbol{g}_{u,n}^* - 1 \right)^2 = 0 \tag{C3}$$

is  $\frac{1}{f_{u,n}^{*}g_{u,n}^{*}}$ . Since  $\tau_n e^{\mu_{u,n}^{*}-1} > 0$ ,  $\tau_n e^{\mu_{u,n}^{*}-1} \left(q_{u,n}^{*}f_{u,n}^{*H}g_{u,n}^{*}-1\right)^2$  is a quadratic function of  $q_{u,n}^{*}$  with an upward opening, and its value first decreases w.r.t.  $q_{u,n}^{*}$  in  $\left[0, \frac{1}{f_{u,n}^{*H}g_{u,n}^{*}}\right]$  and then increases in  $\left[\frac{1}{f_{u,n}^{*H}g_{u,n}^{*}}, +\infty\right)$ . Note that the  $\tau_n e^{\mu_{u'',n}^{*}-1}q_{u,n}^{*2}f_{u'',n}^{*H}g_{u,n}^{*}g_{u,n}^{*H}f_{u'',n}^{*}$  term in (C2) increases w.r.t.  $q_{u,n}^{*}$  in  $\left[0, +\infty\right]$ . Therefore, if  $q_{u,n}^{*} > \frac{1}{f_{u,n}^{*H}g_{u,n}^{*}}$ , we can always decrease  $q_{u,n}^{*}$  such that Eqs. (C1) and (C2) still hold. Then, with the newly obtained  $q_{u,n}^{*}$ , all constraints of (28) hold, and its objective function is further decreased. This contradicts the assumption that  $q_{u,n}^{*}$  is the optimal solution. Therefore, in the optimal case, Eq. (B4) is true. This completes the proof of Lemma B1.