

Related-tweakey impossible differential attack on QARMA-128

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Dear editor,

QARMA [1] is a family of hardware-oriented lightweight tweakable block ciphers designed by Avanzi. It contains two versions, which support block sizes of 64 and 128 bits, denoted by QARMA-64 and QARMA-128, respectively. The structure of QARMA is an Even-Mansour scheme [2] with a keyed pseudo-reflector, and the TWEAKEY framework is taken as an inspiration. Avanzi claimed that QARMA is understood to offer n bits of (time-data tradeoff) security if no better attacks are possible than time $2^{n-d-\epsilon}$ with 2^d chosen or known {plaintext, ciphertext, tweak} triples, for a small ϵ , e.g., 2. The security declaration is under the assumption that the attacker does not have control over the key, but may have full control over the tweak. Impossible differential cryptanalysis was first proposed by Biham et al. [3], the attacker can take advantage of the relation of two tweakeys in related-tweakey scenarios [4]. Yang et al. [5] proposed an impossible differential attack on QARMA, but the results are invalid due to the security claim of the designers.

Differential properties of matrix M . Denote the column-wise input difference of matrix M by ΔIn , and the output difference by ΔOut .

Proposition 1. If $\Delta\text{In} = (0, \alpha, \alpha \ggg 1, \alpha \ggg 4)$, then the output difference $\Delta\text{Out} = (\alpha \lll 3, 0, 0, 0)$, where α is an arbitrary 8-bit value.

The proof and other three propositions are presented in Appendix A.

Distinguishers. Zong et al. [6] proposed an interesting method to search related-key impossible differentials from single key impossible differentials. Inspired by their idea, we also perform automatic search first. However, in order to reduce the data complexity, more active cells at the beginning of the distinguishers are preferred, so that we can construct structures to reduce plaintexts. Therefore, by combining the differential properties proposed above and transition properties [1] of matrix M , we revise the result of our automatic search, and obtain two families of 6-round related-tweakey impossible differential distinguisher for QARMA-128, which are placed between Rounds 7 and 12. We explain distin-

guisher family No.1 in this part and show distinguisher family No.2 in Appendix B.

The entire trail is depicted in Figure 1. The related-tweakey impossible differential characteristic

$$\begin{aligned} &(\alpha(\beta \ggg 4)\varepsilon_1\gamma, \beta(\alpha \ggg 4)(\gamma \ggg 4)\varepsilon_2, \\ &\varepsilon_3(\gamma \ggg 1)(\alpha \ggg 1)(\beta \ggg 1), 0\varepsilon_400) \\ &\rightarrow (0000, \delta 000, 0000, 0000) \end{aligned}$$

exists, if the following conditions are satisfied:

- (1) $\Delta T_7[\gamma]$ is the only active cell of ΔT_7 ;
- (2) $\Delta W_{13}[4]$ is the only active cell of ΔW_{13} , and the equation $\Delta W_{13}[4] = \Delta T_{12}[4] = \delta$ is satisfied, where δ is an arbitrary non-zero value, α, β, γ and ε_i (for $1 \leq i \leq 4$) could be zero or non-zero freely.

Two special cases are explained below, which are adopted in key recovery attacks.

Case 1. If the conditions (1, 2) are satisfied, then the characteristic

$$\begin{aligned} &(00\alpha 0, 000\beta, \gamma 000, 0\sigma 00) \\ &\rightarrow (0000, \delta 000, 0000, 0000) \end{aligned}$$

exists, where $\alpha, \beta, \gamma, \sigma$ and δ are non-zero values.

Case 2. If the conditions (1, 2) are satisfied, then the characteristic

$$\begin{aligned} &(\alpha(\beta \ggg 4)\varepsilon_1\gamma, \beta(\alpha \ggg 4)(\gamma \ggg 4)\varepsilon_2, \\ &\varepsilon_3(\gamma \ggg 1)(\alpha \ggg 1)(\beta \ggg 1), 0\varepsilon_400) \\ &\rightarrow (0000, \delta 000, 0000, 0000) \end{aligned}$$

exists, where $\alpha, \beta, \gamma, \delta$ and ε_i (for $1 \leq i \leq 4$) are non-zero values.

Key recovery attacks. Based on two parallel 6-round distinguishers – distinguisher family No.1 (Case 1) and distinguisher family No.2 (Case 2), we mount a key recovery attack on 10-round QARMA-128 by adding one round before the distinguisher and three rounds after it. The details of this attack are described in Appendix C. The data complexity is $2^{104.02}$ triples while the time complexity is $2^{120.94}$ 10-round encryptions, and the memory complexity is $2^{94.50}$ 128-bit.

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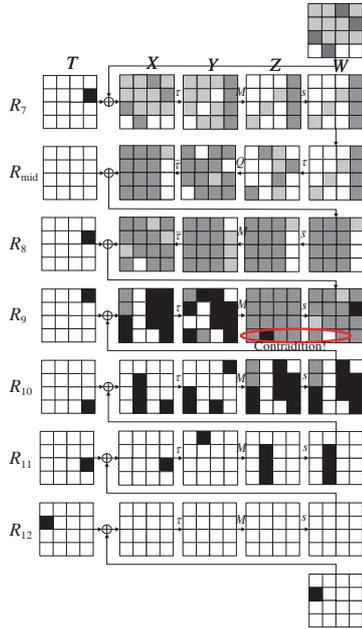


Figure 1 (Color online) Distinguisher family No.1. The black cells are active while the white ones are inactive, the dark and light gray cells could be active or inactive.

We tried to mount a key recovery attack on 11-round QARMA-128 including the outer whitening keys but the product of time and data is greater than 2^{256} . Thus, we proceed an attack on 11-round variant without the outer whitening keys. In order to reduce the time complexity of memory access (offline phase), we derived our attack in a chosen ciphertext-tweak scenario. Based on distinguisher family No.1 (Case 1), the attack requires $2^{102.54}$ triples and $2^{145.98}$ 11-round encryptions. The memory complexity is $2^{135.54}$ 128-bit. The details of this attack are described in Appendix D.

Conclusion. We investigate some interesting properties of the matrix M firstly, which portray the propagation of differential characteristics. Then we construct two families of 6-round related-tweakey impossible differential distinguishers for QARMA-128. Based on two parallel 6-round distinguishers, we proceed key recovery attacks on 10-round

QARMA-128 including the outer whitening keys. We also mount 11-round attacks omitting outer whitening keys in a chosen ciphertext-tweak scenario. Our results are the best impossible differential attacks on QARMA-128 compared to the published ones.

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Supporting information Appendixes A–D. The supporting information is available online at info.scichina.com and link.springer.com. The supporting materials are published as submitted, without typesetting or editing. The responsibility for scientific accuracy and content remains entirely with the authors.

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